#### Clustering: Partition Clustering

### Lecture outline

- Distance/Similarity between data objects
- Data objects as geometric data points
- Clustering problems and algorithms
  - K-means
  - K-median
  - K–center









#### Outliers

• Outliers are objects that do not belong to any cluster or form clusters of very small cardinality



 In some applications we are interested in discovering outliers, not clusters (outlier analysis)

### Why do we cluster?

- Clustering : given a collection of data objects group them so that
  - Similar to one another within the same cluster
  - Dissimilar to the objects in other clusters
- Clustering results are used:
  - As a stand-alone tool to get insight into data distribution
    - Visualization of clusters may unveil important information
  - As a preprocessing step for other algorithms
    - Efficient indexing or compression often relies on clustering

### Applications of clustering?

- Image Processing

   cluster images based on their visual content
- Web
  - Cluster groups of users based on their access patterns on webpages
  - Cluster webpages based on their content
- Bioinformatics
  - Cluster similar proteins together (similarity wrt chemical structure and/or functionality etc)
- Many more...

### The clustering task

 Group observations into groups so that the observations belonging in the same group are similar, whereas observations in different groups are different

#### Basic questions:

- What does "similar" mean
- What is a good partition of the objects? I.e., how is the quality of a solution measured
- How to find a good partition of the observations

# Partitioning algorithms: basic concept

- Construct a partition of a set of n objects into a set of k clusters
  - Each object belongs to exactly one cluster
  - The number of clusters  ${\bf k}$  is given in advance

#### The k-means problem

- Given a set X of n points in a ddimensional space and an integer k
- Task: choose a set of k points { $c_1, c_2, ..., c_k$ } in the d-dimensional space to form clusters { $C_1, C_2, ..., C_k$ } such that  $Cost(C) = \sum_{i=1}^k \sum_{x \in C_i} L_2^2 (x c_i)$

is minimized

• Some special cases: k = 1, k = n

#### Algorithmic properties of the kmeans problem

- NP-hard if the dimensionality of the data is at least 2 (d>=2)
- Finding the best solution in polynomial time is infeasible
- For d=1 the problem is solvable in polynomial time (how?)
- A simple iterative algorithm works quite well in practice

### The k-means algorithm

- One way of solving the k-means problem
- Randomly pick k cluster centers {c<sub>1</sub>,...,c<sub>k</sub>}
- For each i, set the cluster C<sub>i</sub> to be the set of points in X that are closer to c<sub>i</sub> than they are to c<sub>j</sub> for all i≠j
- For each i let c<sub>i</sub> be the center of cluster C<sub>i</sub> (mean of the vectors in C<sub>i</sub>)
- Repeat until convergence

## Properties of the k-means algorithm

- Finds a local optimum
- Converges often quickly (but not always)
- The choice of initial points can have large influence in the result

#### Two different K-means Clusterings



#### Two different K-means Clusterings





#### Two different K-means Clusterings





### Discussion k-means algorithm

- Finds a local optimum
- Converges often quickly (but not always)
- The choice of initial points can have large influence
  - Clusters of different densities
  - Clusters of different sizes
- Outliers can also cause a problem (Example?)

# Some alternatives to random initialization of the central points

- Multiple runs
  - Helps, but probability is not on your side
- Select original set of points by methods other than random . E.g., pick the most distant (from each other) points as cluster centers (kmeans++ algorithm)

### The k-median problem

- Given a set X of n points in a ddimensional space and an integer k
- Task: choose a set of k points {c<sub>1</sub>,c<sub>2</sub>,...,c<sub>k</sub>} from X and form clusters {C<sub>1</sub>,C<sub>2</sub>,...,C<sub>k</sub>} such that

$$Cost(C) = \sum_{i=1}^{k} \sum_{x \in C_i} L_1(x, c_i)$$

#### is minimized

#### The k-medoids algorithm

- Or ... PAM (Partitioning Around Medoids, 1987)
  - Choose randomly k medoids from the original dataset X
  - Assign each of the n-k remaining points in X to their closest medoid
  - iteratively replace one of the medoids by one of the non-medoids if it improves the total clustering cost

#### Discussion of PAM algorithm

 The algorithm is very similar to the kmeans algorithm

It has the same advantages and disadvantages

• How about efficiency?

# CLARA (Clustering Large Applications)

- It draws **multiple samples** of the data set, applies PAM on each sample, and gives the best clustering as the output
- <u>Strength</u>: deals with larger data sets than PAM
- <u>Weakness:</u>
  - Efficiency depends on the sample size
  - A good clustering based on samples will not necessarily represent a good clustering of the whole data set if the sample is biased

### The k-center problem

- Given a set X of n points in a ddimensional space and an integer k
- Task: Partition the points in X in k clusters {C<sub>1</sub>,C<sub>2</sub>,...,C<sub>k</sub>} such that

$$R(C) = \max_{i} \max_{x, x' \in C_i} d(x, x')$$

is minimized

#### Algorithmic properties of the kcenters problem

- NP-hard if the dimensionality of the data is at least 2 (d>=2)
- Finding the best solution in polynomial time is infeasible
- For d=1 the problem is solvable in polynomial time (how?)
- A simple combinatorial algorithm works well in practice

# The furthest-first traversal algorithm

- Pick any data point and label it as point 1
- For i=2,3,...,k
  - Find the unlabelled point furthest from {1,2, ...,i-1} and label it as i.
    - //Use  $d(x,S) = \min_{y \in S} d(x,y)$  to identify the

distance //of a point from a set

 $-\pi(i) = \operatorname{argmin}_{j < i} d(i,j)$ 

 $- R_i = d(i, \pi(i))$ 

• Assign the remaining unlabelled points to their closest labelled point

The furthest-first traversal is a 2-approximation algorithm

• Claim1:  $R_1 \ge R_2 \ge ... \ge R_n$ 

• Proof:  $-R_{j}=d(j,\pi(j)) = d(j,\{1,2,...,j-1\})$   $\leq d(j,\{1,2,...,i-1\}) //j > i$   $\leq d(i,\{1,2,...,i-1\}) = R_{i}$ 

## The furthest-first traversal is a 2-approximation algorithm

 Claim 2: If C is the clustering reported by the farthest algorithm, then
 R(C)=R<sub>k+1</sub>

#### • Proof:

#### - For all i > k we have that $d(i, \{1,2,...,k\}) \le d(k+1,\{1,2,...,k\}) = R_{k+1}$

# The furthest-first traversal is a 2-approximation algorithm

- Theorem: If C is the clustering reported by the farthest algorithm, and C<sup>\*</sup>is the optimal clustering, then R(C)≤2xR(C<sup>\*</sup>)
- Proof:
  - Let C\*1, C\*2,..., C\*k be the clusters of the optimal kclustering.
  - If these clusters contain points  $\{1,...,k\}$  then  $R(C) \le 2R(C^*)$  (triangle inequality)
  - Otherwise suppose that one of these clusters contains two or more of the points in {1,...,k}. These points are at distance at least R<sub>k</sub> from each other. Thus clusters must have radius  $\frac{1}{2}$  R<sub>k</sub>  $\geq \frac{1}{2}$  R<sub>k+1</sub> =  $\frac{1}{2}$  R(C)

## What is the right number of clusters?

- ... or who sets the value of k?
- For n points to be clustered consider the case where k=n. What is the value of the error function
- What happens when  $\mathbf{k} = \mathbf{1}$ ?
- Since we want to minimize the error why don't we select always k = n?

# Occam's razor and the minimum description length principle

- Clustering provides a description of the data
- For a description to be good it has to be:
  - Not too general
  - Not too specific
- Penalize for every extra parameter that one has to pay
- Penalize the number of bits you need to describe the extra parameter
- So for a clustering C, extend the cost function as follows:
- NewCost(C) = Cost(C) +  $|C| \times logn$