

# Advanced Topics in Data Mining

## Special focus: Social Networks

# Reminders

- By the end of this week/ beginning of next we need to have a tentative presentation schedule
- Each one of you should send me an email about a theme by Friday, February 22.

What did we learn in the last lecture?

# What did we learn in the last lecture?

- Degree distribution
  - What are the observed degree distributions
- Clustering coefficient
  - What are the observed clustering coefficients?
- Average path length
  - What are the observed average path lengths?

# What are we going to learn in this lecture?

- How to generate graphs that have the desired properties
  - Degree distribution
  - Clustering coefficient
  - Average path length
- We are going to talk about **generative models**

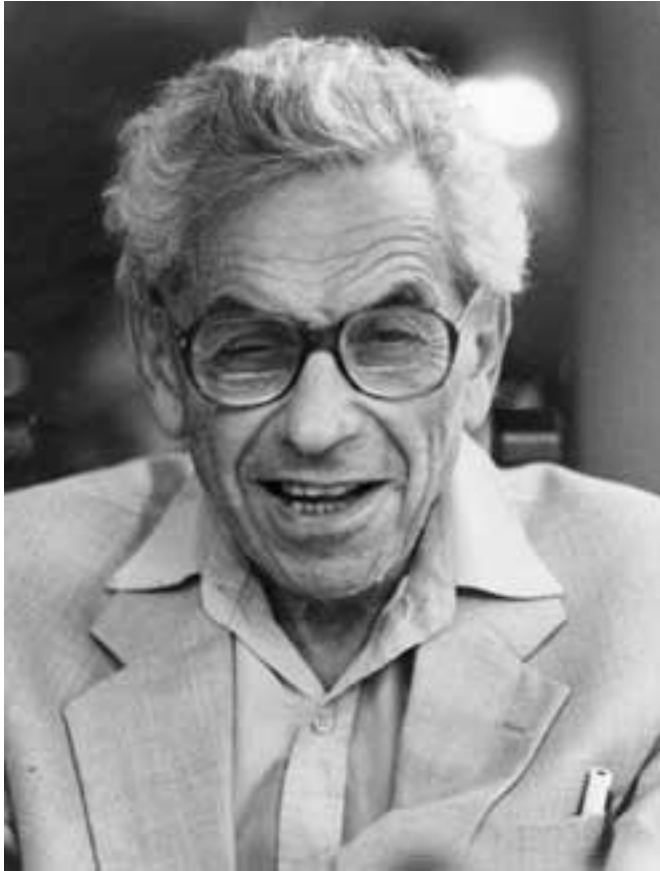
# What is a network model?

- Informally, a network model is a **process** (radomized or deterministic) for generating a graph
- Models of **static** graphs
  - **input**: a set of parameters  $\Pi$ , and the size of the graph  $n$
  - **output**: a graph  $G(\Pi, n)$
- Models of **evolving** graphs
  - **input**: a set of parameters  $\Pi$ , and an initial graph  $G_0$
  - **output**: a graph  $G_t$  for each time  $t$

# Families of random graphs

- A deterministic model  $D$  defines a single graph for each value of  $n$  (or  $t$ )
- A randomized model  $R$  defines a probability space  $\langle G_n, P \rangle$  where  $G_n$  is the set of all graphs of size  $n$ , and  $P$  a probability distribution over the set  $G_n$  (similarly for  $t$ )
  - we call this a family of random graphs  $R$ , or a random graph  $R$

# Erdős-Renyi Random graphs



Paul Erdős (1913-1996)



# Erdős-Renyi Random Graphs

- The  $G_{n,p}$  model
  - **input**: the number of vertices  $n$ , and a parameter  $p$ ,  $0 \leq p \leq 1$
  - **process**: for each pair  $(i,j)$ , generate the edge  $(i,j)$  independently with probability  $p$
- Related, but not identical: The  $G_{n,m}$  model
  - **process**: select  $m$  edges uniformly at random

# Graph properties

- A property  $P$  holds **almost surely** (or for **almost every** graph), if

$$\lim_{n \rightarrow \infty} P[G \text{ has } P] = 1$$

- Evolution of the graph: which properties hold as the probability  $p$  increases?
- **Threshold phenomena**: Many properties appear suddenly. That is, there exist a probability  $p_c$  such that for  $p < p_c$  the property does not hold a.s. and for  $p > p_c$  the property holds a.s.
  - ***What do you expect to be a threshold phenomenon in random graphs?***

# The giant component

- Let  $z=np$  be the average degree
- If  $z < 1$ , then almost surely, the largest component has size at most  $O(\ln n)$
- if  $z > 1$ , then almost surely, the largest component has size  $\Theta(n)$ . The second largest component has size  $O(\ln n)$
- if  $z = \omega(\ln n)$ , then the graph is almost surely connected.

# The phase transition

- When  $z=1$ , there is a phase transition
  - The largest component is  $O(n^{2/3})$
  - The sizes of the components follow a power-law distribution.

# Random graphs degree distributions

- The degree distribution follows a **binomial**

$$p(k) = B(n; k; p) = \binom{n}{k} p^k (1-p)^{n-k}$$

- Assuming  $z=np$  is fixed, as  $n \rightarrow \infty$ ,  $B(n, k, p)$  is approximated by a **Poisson** distribution

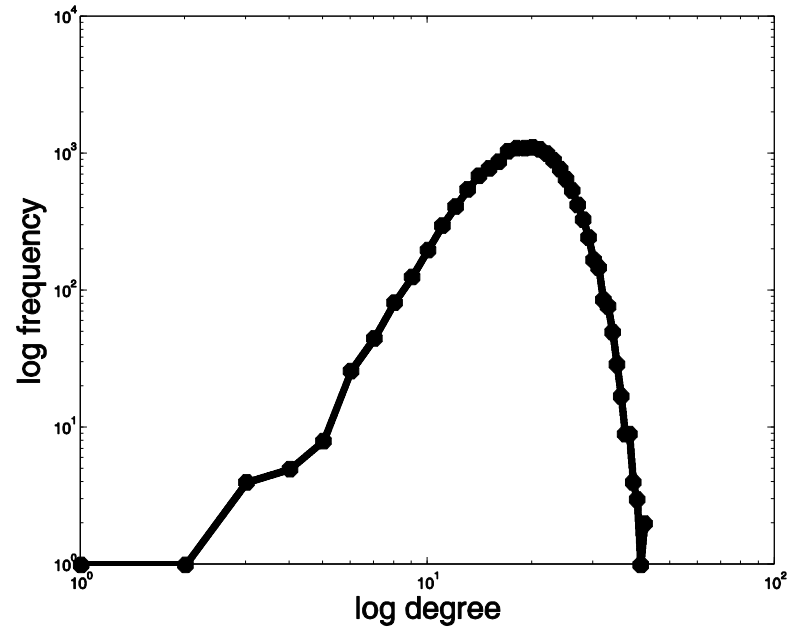
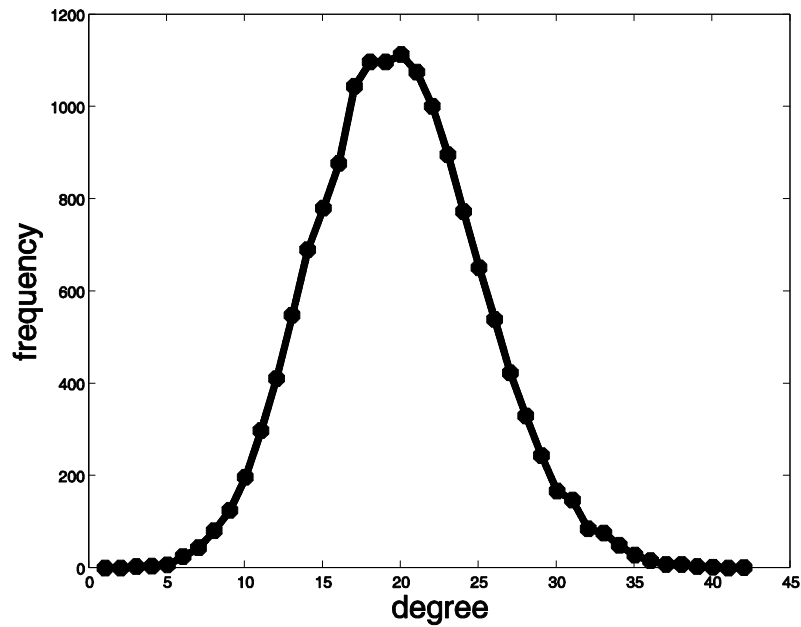
$$p(k) = P(k; z) = \frac{z^k}{k!} e^{-z}$$

- Highly concentrated around the mean, with a tail that drops exponentially

# Random graphs and real life

- A beautiful and elegant theory studied exhaustively
- Random graphs had been used as idealized network models
- Unfortunately, they don't capture reality...

# A random graph example



# Departing from the Random Graph model

- We need models that better capture the characteristics of real graphs
  - degree sequences
  - clustering coefficient
  - short paths



# Graphs with given degree sequences

- **input:** the degree sequence  $[d_1, d_2, \dots, d_n]$
- Can you generate a graph with nodes that have degrees  $[d_1, d_2, \dots, d_n]$  ?
- ? 😊

# Graphs with given degree sequences

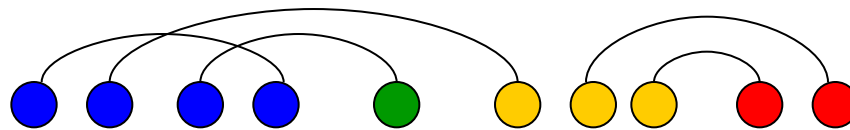
- The configuration model
  - **input:** the degree sequence  $[d_1, d_2, \dots, d_n]$
  - **process:**
    - Create  $d_i$  copies of node  $i$
    - Take a random matching (pairing) of the copies
      - self-loops and multiple edges are allowed
- Uniform distribution over the graphs with the given degree sequence

# Example

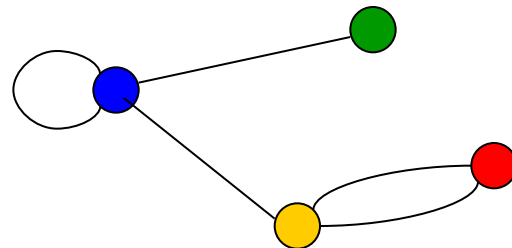
- Suppose that the degree sequence is



- Create multiple copies of the nodes



- Pair the nodes uniformly at random
- Generate the resulting network



# Graphs with given degree sequences

- How about **simple** graphs ?
  - No self loops
  - No multiple edges

# Graphs with given degree sequences

- Realizability of degree sequences
- **Lemma:** A degree sequence  $\mathbf{d} = [d(1), \dots, d(n)]$  with  $d(1) \geq d(2) \geq \dots \geq d(n)$  and  $d(1) + d(2) + \dots + d(n)$  **even** is **realizable** if and only if for every  $1 \leq k \leq n-1$  it holds that

$$\sum_{i=1}^k d(i) \leq k(k-1) + \sum_{i=k+1}^n \min\{k, d(i)\}$$

# Graphs with given degree sequences -- algorithm

- Input :  $\mathbf{d} = [d(1), \dots, d(n)]$
- Output: No or simple graph  $\mathbf{G} = (\mathbf{V}, \mathbf{E})$  with degree sequence  $\mathbf{d}$
- If  $\sum_{i=1 \dots n} d(i)$  is odd return "No"
- While 1 do
  - If there exist  $i$  with  $d(i) < 0$  return "No"
  - If  $d(i) = 0$  for all  $i$  return the graph  $\mathbf{G} = (\mathbf{V}, \mathbf{E})$
  - Pick random node  $\mathbf{v}$  with  $d(\mathbf{v}) > 0$
  - $\mathbf{S}(\mathbf{v}) =$  set of nodes with the  $d(\mathbf{v})$  highest  $\mathbf{d}$  values
  - $d(\mathbf{v}) = 0$
  - For each node  $\mathbf{w}$  in  $\mathbf{S}(\mathbf{v})$ 
    - $\mathbf{E} = \mathbf{E} \setminus \text{union}(\mathbf{v}, \mathbf{w})$
    - $d(\mathbf{w}) = d(\mathbf{w}) - 1$

How can we generate data with power-law degree distributions?

# Preferential Attachment in Networks

- First considered by [Price 65] as a model for citation networks
  - each new paper is generated with  $m$  citations (mean)
  - new papers cite previous papers with probability proportional to their indegree (citations)
  - what about papers without any citations?
    - each paper is considered to have a “default” citation
    - probability of citing a paper with degree  $k$ , proportional to  $k+1$
- Power law with exponent  $\alpha = 2+1/m$



# Barabasi-Albert model

- The BA model (undirected graph)
  - **input**: some initial subgraph  $G_0$ , and  $m$  the number of edges per new node
  - **the process**:
    - nodes arrive one at the time
    - each node connects to  $m$  other nodes selecting them with probability proportional to their degree
    - if  $[d_1, \dots, d_t]$  is the degree sequence at time  $t$ , the node  $t+1$  links to node  $i$  with probability

$$\frac{d_i}{\sum_i d_i} = \frac{d_i}{2mt}$$

- Results in power-law with exponent  $\alpha = 3$

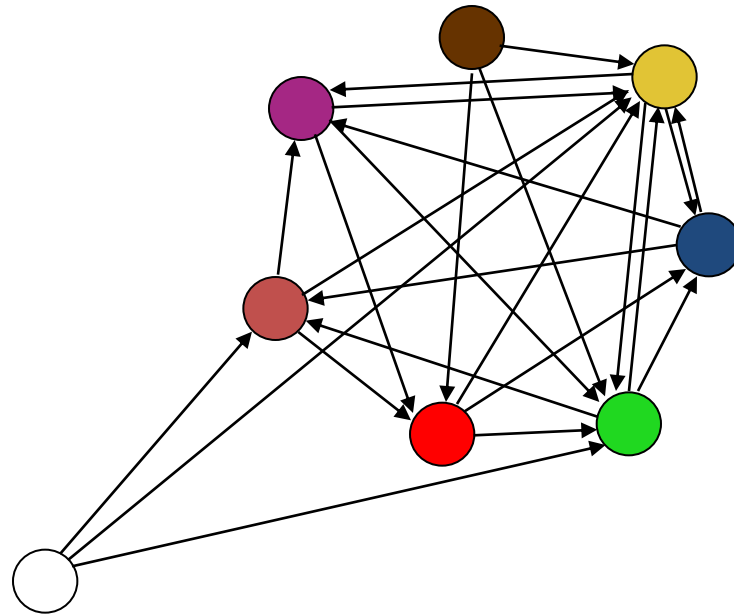
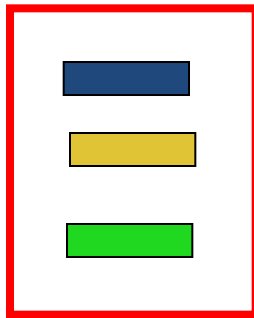
# Variations of the BA model

- Many variations have been considered

# Copying model

- Input:
  - the out-degree  $d$  (constant) of each node
  - a parameter  $\alpha$
- The process:
  - Nodes arrive one at the time
  - A new node selects uniformly one of the existing nodes as a **prototype**
  - The new node creates  $d$  outgoing links. For the  $i^{\text{th}}$  link
    - with probability  $\alpha$  it copies the  $i$ -th link of the prototype node
    - with probability  $1 - \alpha$  it selects the target of the link uniformly at random

# An example



# Copying model properties

- Power law degree distribution with exponent  $\beta = (2-\alpha)/(1-\alpha)$
- Number of bipartite cliques of size  $i \times d$  is  $ne^{-i}$
- The model has also found applications in biological networks
  - copying mechanism in gene mutations

# Small world Phenomena

- So far we focused on obtaining graphs with power-law distributions on the degrees. What about other properties?
  - **Clustering coefficient**: real-life networks tend to have high clustering coefficient
  - **Short paths**: real-life networks are “**small worlds**”
    - this property is easy to generate
  - Can we combine these two properties?

# Small-world Graphs

- According to Watts [W99]
  - Large networks ( $n \gg 1$ )
  - Sparse connectivity (avg degree  $z \ll n$ )
  - No central node ( $k_{\max} \ll n$ )
  - Large clustering coefficient (larger than in random graphs of same size)
  - Short average paths ( $\sim \log n$ , close to those of random graphs of the same size)

# Mixing order with randomness

- Inspired by the work of Solmonoff and Rapoport
  - nodes that share neighbors should have higher probability to be connected
- Generate an edge between  $i$  and  $j$  with probability proportional to  $R_{ij}$

$$R_{ij} = \begin{cases} 1 & \text{if } m_{ij} \geq z \\ \left(\frac{m_{ij}}{z}\right)^\alpha (1-p) + p & \text{if } 0 < m_{ij} < z \\ p & \text{if } m_{ij} = 0 \end{cases}$$

$m_{ij}$  = number of common neighbors of  $i$  and  $j$

$p$  = very small probability

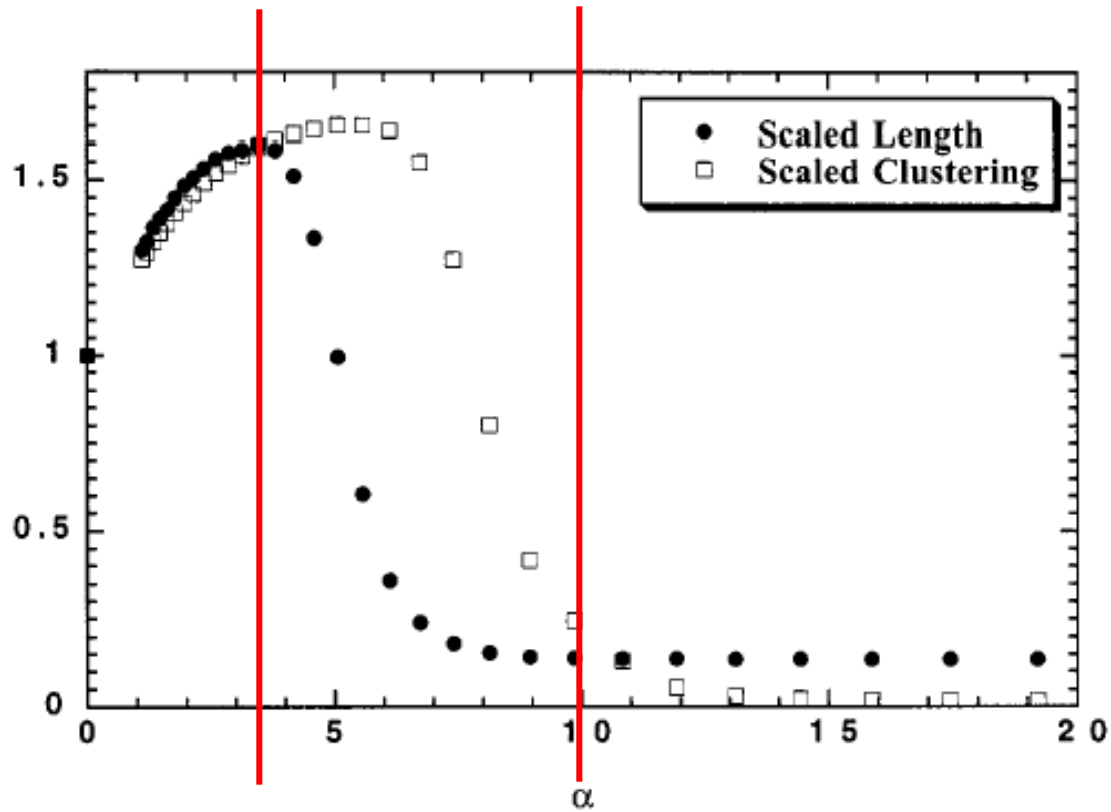
- When  $\alpha = 0$ , edges are determined by common neighbors
- When  $\alpha = \infty$  edges are independent of common neighbors
- For intermediate values we obtain a combination of order and randomness



# Algorithm

- Start with a ring
- For  $i = 1 \dots n$ 
  - Select a vertex  $j$  with probability proportional to  $R_{ij}$  and generate an edge  $(i,j)$
- Repeat until  $z$  edges are added to each vertex

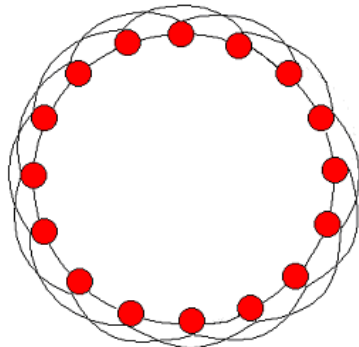
# Clustering coefficient – Avg path length



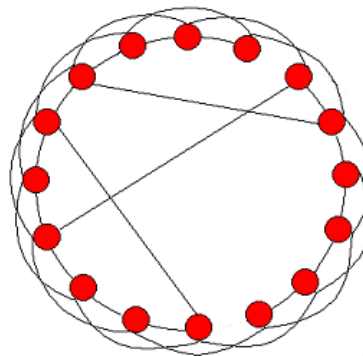
small world graphs

# Watts and Strogatz model [WS98]

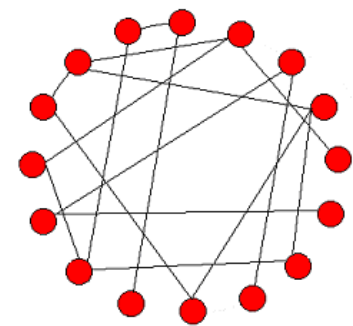
- Start with a ring, where every node is connected to the next  $z$  nodes
- With probability  $p$ , **rewire** every edge (or, add a **shortcut**) to a uniformly chosen destination.
  - Granovetter, “The strength of weak ties”



order



$0 < p < 1$



randomness

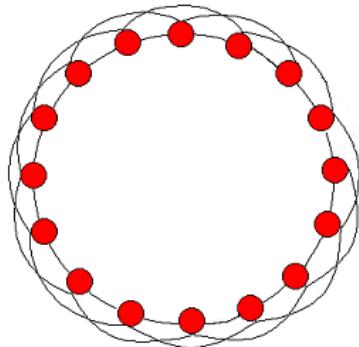
$p = 0$

$p = 1$

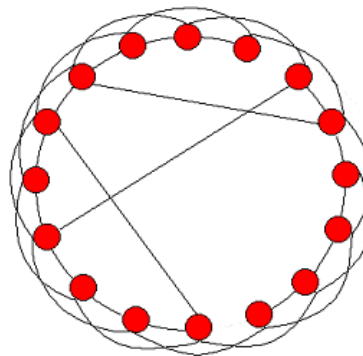


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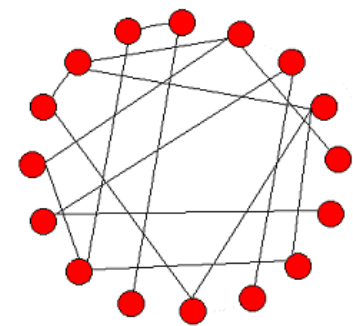
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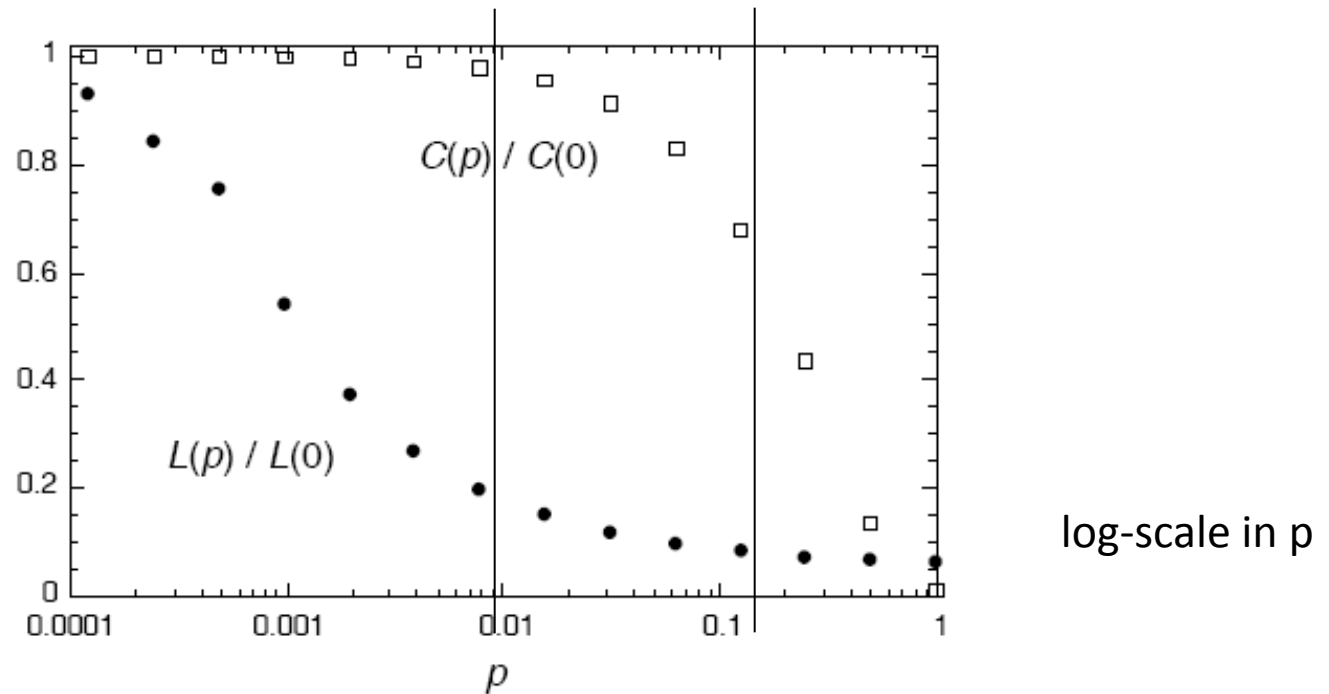
randomness

$p = 0$

$p = 1$



# Clustering Coefficient – Characteristic Path Length



When  $p = 0$ ,  $C = 3(k-2)/4(k-1) \sim 3/4$   
 $L = n/k$

For small  $p$ ,  $C \sim 3/4$   
 $L \sim \log n$

# Next Class

- Some more generative models for social-network graphs