Reconstructing randomized graphs

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Objective

 Reconstruct a graph and/or feature set that has been randomized by a known algorithm

- Why
 - Attack privacy
 - Lossy/Noisy information
 - Data analysis (Congressional voting?)

Problem Definition

- Given an observed graph (G') and feature set (F') try to rebuild the original G and F.
- Possible observed pairs:
 - G' = G and F' != F (G is not randomized, F is)
 - G' != G and F' = F (F is not randomized, G is)
 - G' != G and F' != F (both G and F are randomized)
- All possible observations can be solved in Polynomial time!!!!

- Graph (G) represented by an adjacency matrix
 - seen this before 1 represents edge, 0 no edge
 - $g_{_{II}}$ represents edge in G
- Feature set (F) represented by matrix of 0 1 values
 - row represents a node, column represents a feature
 - if graph has n nodes and each node has k features there are nk entries in matrix
 - \mathbf{f}_{i} is a feature vector in F, \mathbf{f}_{i} represents feature

- Relationship between nodes and features?
 - From Plato, "Friends have all things in common."
 - Key assumption, relates features and nodes
- Edge exists if two nodes have features in common
 - Use a similarity function between feature vectors

- Similarity Function types used in the paper:
 - Dot Product (DP): $sim(\mathbf{f}_i, \mathbf{f}_j) = \sum_{l=1}^{n} f_{il} \cdot f_{jl}$

• Hamming (H):
$$sim(f_i, f_j) = \sum_{l=1}^{k} 1 - |f_{il} - f_{jl}|$$

• More similar features, more probable an edge

$$Pr(g_{ij}=1|f_i,f_j) = \frac{1}{Z}e^{\alpha \sin(f_i,f_j)}$$
$$Pr(g_{ij}=0|f_i,f_j) = \frac{1}{Z}e^{\alpha(1-\sin(f_i,f_j))}$$

• Assume edges are independent and features define edges

$$Pr(G,F) = Pr(G|F) = \prod_{i < j} Pr(g_{ij}|f_i,f_j)$$

- Randomization From matrix X produce X'
 - Key probabilities needed:
 - Pr(x=0|x=0) and Pr(x=1|x=0)
 - Pr(x=0|x=1) and Pr(x=1|x=1)
 - Must be able to calculate Pr(X|X')

$$Pr(X|X') = \frac{Pr(X'|X)Pr(X)}{Pr(X')} \propto Pr(X'|X) = \prod_{i,j} Pr(x'_{ij}|x_{ij})$$

Reconstruction

- Uses a maximum likelihood approach
 - Given observation which actual is most probable?
- Three types of problems:
 - G reconstruction (when F not altered)
 - F reconstruction (when G not altered)
 - GF reconstruction (when both altered)
- In short want to find Pr(G,F | G',F')

Reconstruction

- Instead of maximizing Pr(G, F | G', F'), minimize the -log(Pr(G,F | G', F'))
- How?

$$Pr(G,F|G',F') = \frac{Pr(G',F'|G,F)Pr(G,F)}{Pr(G',F')}$$

 $Pr(G,F|G',F') \propto Pr(G',F'|G,F) Pr(G,F) = Pr(G'|G) Pr(F'|F) Pr(G,F)$

Energy Function for minimization -

 $E(G,F) = -\log Pr(G,F|G'F') = -\log Pr(G'|G) - \log Pr(F'|F) - \log Pr(G,F)$

Reconstruction

• GF – reconstruction energy function:

 $E(G,F) = -\log Pr(G,F|G'F') = -\log Pr(G'|G) - \log Pr(F'|F) - \log Pr(G,F)$

• G – reconstruction energy function:

 $E(G) = -\log Pr(G'|G) - \log Pr(G,F) = \sum_{i < j} (-\log Pr(g'_{ij}|g_{ij}) - \log Pr(g_{ij}|f'_i,f'_j))$

• F – reconstruction energy function: $E(F) = -\log Pr(F'|F) - \log Pr(G,F) = \sum_{i=1}^{n} \sum_{l=1}^{k} (-\log Pr(f'_{il}|f_{il}) - \log Pr(g'_{ij}|f_i,f_j))$

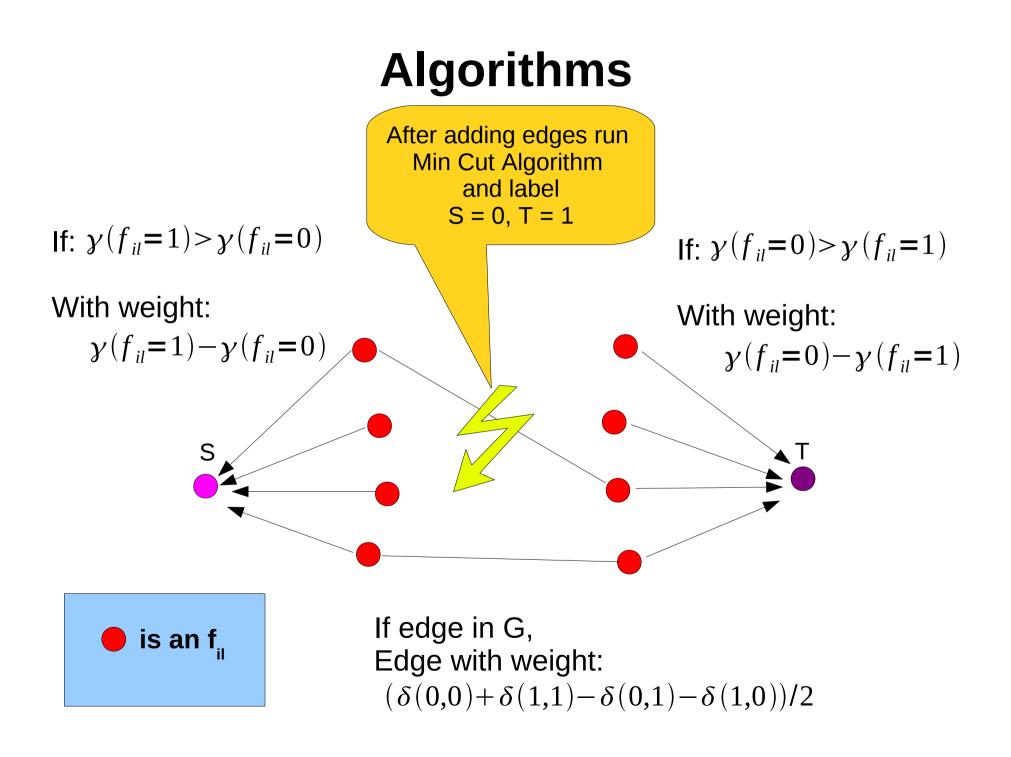
- G reconstruction, rebuild edges
 - For every two nodes, calculate the energy of an edge using E(G).
 - If E(edge) < E(no edge) then add an edge else no edge, remember we are trying to minimize energy
 - Optimal algorithm but not guaranteed to rebuild the original graph
 - Running time: $O(T_s n^2)$

- F reconstruction, label feature values 0-1
 - Optimal algorithm from computer vision
 - Uses Min cut algorithm in a unique way (very cool!!)
 - Polynomial solution but expensive for computation time and space requirements
 - Naïve suboptimal algorithm
 - Performs labeling in greedy fashion
 - Makes assignment that best minimizes E in that move
 - Iteration based...

- Optimal F reconstruction
 - Intuition, assign labels (1/0) to nk features
 - First rewrite E(F):

$$E(F) = \sum_{i=1}^{n} \sum_{l=1}^{k} \left(-\log \Pr(f'_{il} | f_{il}) - \log \Pr(g'_{ij} | f_i, f_j) \right) = \sum_{i=1}^{n} \sum_{l=1}^{k} \left(\gamma(f_{il}) - \delta(f_i, f_j) \right)$$

 Next build a flow-graph where each f_i is a node v_i and add two terminals s, t.



Connection to terminals

- To s if $\gamma(f_{il}=1) > \gamma(f_{il}=0)$
 - with weight $\gamma(f_{il}=1)-\gamma(f_{il}=0)$
- Totif $\gamma(f_{il}=0) > \gamma(f_{il}=1)$
 - with weight $\gamma(f_{il}=0)-\gamma(f_{il}=1)$

Connection to other nodes

- When edge in G, add edge to flow graph
 - with weight $(\delta(0,0)+\delta(1,1)-\delta(0,1)-\delta(1,0))/2$

- Perform min cut of flow graph
 - Nodes attached to s are labeled 0
 - Nodes attached to t are labeled 1
- Pretty Cool? Huh?

• More Info V. Komogorov, R. Zabih What Energy Functions Can Be Minimized via Graph Cuts?

- GF reconstruction
 - Similar to F reconstruction
 - Naïve Algorithm
 - Assigns values that minimize energy based on current move; greedy, suboptimal
 - Optimal Algorithm
 - Same as F-reconstruction but with a few modifications

- Optimal GF reconstruction
 - Intuition assign labels (1/0) to nk feature nodes, nC2 edges, and k(nC2) triples representing edge feature relationship
 - Requires manipulations of E(G,F), not presented in paper but explained as 'simple'
 - Limited to DP similarity functions only
 - Restriction based on behavior of energy function

- Manipulations result in three new edge evaluations
 - For \mathbf{g}_{ij} : $\sigma_g(g_{ij}) = \alpha k g_{ij} \log Pr(g'_{ij}|g_{ij})$

• For
$$\mathbf{f}_{il}$$
: $\sigma_f(f_{il}) = -\log Pr(f'_{il}|f_{il})$

- For \mathbf{g}_{ij} , \mathbf{f}_{ij} , \mathbf{f}_{jj} tuples: $\sigma(g_{ij}, f_{il}, f_{jl}) = \alpha(1-2g_{ij}) \sin(f_{il}, f_{jl})$
- Evaluations of edges occurs identically to the optimal F reconstruction algorithm

- Computational speedups
 - Optimal F and GF and naïve are very expensive
 - Solution is to divide the input space up, solve each subdivision, and aggregate results
 - Proposed algorithm uses a BFS tree

- Tested algorithms against 3 datasets
 - A synthetic built dataset with n=200 and k=20
 - Controlled edge probability for 557 edges
 - DBLP dataset of author publications n=4981, k=19, and 20670 edges
 - Terror dataset
 - Nodes are attacks, features are attack characteristics, and edge exists if attack occurred at same location
 - N=645, k=94, and edges=3172

- G reconstruction
 - Used DBLP dataset
 - Tested against data subjected to increasing randomization amounts
 - Optimal algorithm performs pretty well in the presence of noisy data.
 - Error rate stabilizes at .625 as randomization increases

- F reconstruction
 - Used synthetic dataset
 - Tested against data subjected to increasing randomization amounts
 - Bounded naïve algorithm iterations to "clock time" of optimal solution
 - For small amounts of randomization naïve and optimal are close in error rate, for larger levels naïve performs worse than optimal

- GF reconstruction
 - Used all 3 datasets, only reported findings of Terror
 - Tested against data subjected to increasing randomization amounts
 - Various results:
 - For DP sim: OptBoth with split better than naïve both and split naïve both
 - For H sim: naïve both better than naïve both with split

 DP sim function "lures" the reconstruction methods to fill up entries with 1s. Can corrected for by proper tuning of edge probability function

 Also noted that the objective function might result in (G,F) with high likelihood but low structural similarity to the data

Project Details

- In General
 - Rebuilding friendships from group information on Facebook
 - Are groups enough on Facebook to define friendships? Do we need more?
 - Challenges are getting a comprehensive dataset with enough group information for analysis
 - Thoughts?

The End

• Questions?

• Thank you