Comp115: Databases

SQL: The Query Language

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Today's course

intuitive way to ask queries

unlike procedural languages (C/C++, java)

[which specify how to solve a problem (or answer a question)]

SQL is a declarative query language

[we ask what we want and the DBMS is going to deliver]

Introduction to SQL

SQL is a relational **query language** supports **simple** yet **powerful** *querying* of data It has two parts:

DDL: Data Definition Language (define and modify schema) (we discussed about that in Relational Model)

DML: Data Manipulation Language (intuitively query data)

Reiterate some terminology

Relation (or table)

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@ee	18	3.2

Row (or tuple)

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@ee	18	3.2

Column (or attribute)

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@ee	18	3.2

Reiterate some terminology

Primary Key (PK)

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@ee	18	3.2

The PK of a relation is the column (or the group of columns) that can uniquely define a row.

In other words:

Two rows **cannot** have the same PK.

The simplest SQL query

"find all contents of a table" in this example: "Find all info for all students"

SELECT *
FROM Students S

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@ee	18	3.2
53777	White	white@cs	19	4.0

to find just names and logins, replace the first line:

SELECT S.name, S.login

Show specific columns

"find name and login for all students"

SELECT S.name, S.login FROM Students S

name	login
Jones	jones@cs
Smith	smith@ee
White	white@cs

this is called: "project name and login from table Students"

Show specific rows

"find all 18 year old students"

SELECT *
FROM Students S
WHERE S.age=18

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@ee	18	3.2

this is called: "select students with age 18."

Querying Multiple Relations

can specify a join over two tables as follows:

SELECT S.name, E.cid FROM Students S, Enrolled E WHERE S.sid=E.sid AND E.grade='B'

sid	cid	grade	
53831	Carnatic101	C	
53831	Reggae203	В	
53650	Topology112	A	
53666	History 105	В	

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@ee	18	3.2

result =

S.name	E.cid
Jones	History105

Basic SQL Query FROM

[DISTINCT] target-list relation-list qualification

relation-list: a list of relation names

target-list: a list of attributes of tables in relation-list

qualification: comparisons using AND, OR and NOT

comparisons are: $\langle op \rangle \langle const \rangle$ or $\langle attr1 \rangle \langle op \rangle \langle attr2 \rangle$, where op is:

$$<,>,=,\leq,\geq,\neq$$

DISTINCT: optional, removes duplicates

By default SQL SELECT does *not* eliminate duplicated! (result is called a "multiset")

Query Semantics

Conceptually, a SQL query can be computed:

- (1) **FROM**: compute <u>cross-product</u> of tables (e.g., Students and Enrolled)
- (2) **WHERE**: Check conditions, discard tuples that fail (applying "selection" condition)
 - (3) **SELECT**: Delete unwanted fields (applying "projection")
 - (4) if **DISTINCT** specified, eliminate duplicate rows

probably the least efficient way to compute a query! **Query Optimization** finds the *same answer* more efficiently

Remember the query and the data

sid	cid	grade
53831	Carnatic101	C
53831	Reggae203	В
53650	Topology112	A
53666	History105	В

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@ee	18	3.2

Step 1 – Cross Product

Combine with cross-product all tables of the **FROM** clause.

S.sid	S.name	S.login	S.age	S.gpa	E.sid	E.cid	E.grade
53666	Jones	jones@cs	18	3.4	53831	Carnatic101	C
53666	Jones	jones@cs	18	3.4	53832	Reggae203	В
53666	Jones	jones@cs	18	3.4	53650	Topology112	A
53666	Jones	jones@cs	18	3.4	53666	History105	В
53688	Smith	smith@ee	18	3.2	53831	Carnatic101	C
53688	Smith	smith@ee	18	3.2	53831	Reggae203	В
53688	Smith	smith@ee	18	3.2	53650	Topology112	A
53688	Smith	smith@ee	18	3.2	53666	History105	В

Step 2 - Discard tuples that fail predicate

Make sure the **WHERE** clause is true!

_	S.sid	S.name	S.login	S.age	S.gpa	E.sid	E.cid	E.grade
	53666	Jones	jones@cs	18	3.4	53831	Carnatic101	C
	53666	Jones	jones@cs	18	3.4	53832	Reggae203	(B)
	53666	Jones	jones@cs	18	3.4	53650	Topology112	Ā
	53666		jones@cs		3.4		History 105	B
	53688	Smith	smith@ee	18	3.2	53831	Carnatic 101	C
	53688	Smith	smith@ee	18	3.2	53831	Reggae203	(B)
	53688	Smith	smith@ee	18	3.2	53650	Topology112	A
	53688	Smith	smith@ee	18	3.2	53666	History105	B

Step 3 - Discard Unwanted Columns

Show only what is on the **SELECT** clause.

S.sid	S.name	S.login	S.age	S.gpa	E.sid	E.cid	E.grade
53666	Jones	jones@cs	18	3.4	53831	Carnatic101	S
53666	Jones	jones@cs	18	3.4	53832	Reggae203	B
53666	Jones	jones@cs	18	3.4	53650	Topology112	A
53666	Jones	jones@cs	18	3.4	53666	History105	(B)
53688	Smith	smith@ee	18	3.2	53831	Carnatic101	C
53688	Smith	smith@ee	18	3.2	53831	Reggae203	(B)
53688	Smith	smith@ee	18	3.2	53650	Topology112	A
53688	Smith	smith@ee	18	3.2	53666	History105	(B)

Reserves

Now the Details...

We will use these instances of relations in our examples. *Sailors*

sid	<u>bid</u>	<u>day</u>
22	101	10/10/16
95	103	11/12/16

sid	sname	rating	age
22	Dustin	7	45.0
31	Lubber	8	55.5
95	Bob	3	63.5

Boats

<u>bid</u>	bname	color
101	Interlake	blue
102	Interlake	red
103	Clipper	green
104	Marine	red

Another Join Query

SELECT sname

FROM Sailors, Reserves

WHERE Sailors.sid=Reserves.sid

AND bid=103

(sid)	sname	rating	age	(sid)	bid	day
22	dustin	7	45.0	22	101	10/10/16
22	dustin	7	45.0	95	103	11/12/16
31	lubber	8	55.5	22	101	10/10/16
31	lubber	8	55.5	95	103	11/12/16
95	Bob	3	63.5	22	101	10/10/16
95	Bob	3	63.5	95	103	11/12/16

Range Variables

can associate "range variables" with the tables in the FROM clause

```
a shorthand, like the <u>rename operator</u> from last time saves writing, makes queries easier to understand "FROM Sailors, Reserves"

"FROM Sailors S, Reserves R"

needed when ambiguity could arise for example, if same table used multiple times in same FROM (called a "self-join")

"FROM Sailors S1, Sailors S2"
```

Range Variables

```
SELECT sname FROM Sailors, Reserves WHERE Sailors.sid=Reserves.sid AND bid=103
```

can be rewritten using range variables as:

SELECT S.sname FROM Sailors S, Reserves R WHERE S.sid=R.sid AND bid=103

Range Variables

an example requiring range variables (self-join)

```
SELECT S1.sname, S1.age, S2.sname, S2.age
FROM Sailors S1, Sailors S2
WHERE S1.age > S2.age
```

another one: "*" if you don't want a projection:

```
SELECT *
FROM Sailors S
WHERE S.age > 20
```

Find sailors who've reserved at least one boat

SELECT S.sid FROM Sailors S, Reserves R WHERE S.sid=R.sid

does DISTINCT makes a difference?



what is the effect of replacing *S.sid* by *S.sname* in the SELECT clause?

Would adding DISTINCT to this variant of the query make a difference?

Expressions

Can use arithmetic expressions in SELECT clause (plus other operations we'll discuss later)

Use AS to provide column names

```
SELECT S.age, S.age-5 AS age1, 2*S.age AS age2 FROM Sailors S
WHERE S.sname = 'dustin'
```

Can also have expressions in WHERE clause:

```
SELECT S1.sname AS name1, S2.sname AS name2 FROM Sailors S1, Sailors S2 WHERE 2*S1.rating = S2.rating - 1
```

String operations

SQL also supports some string operations "LIKE" is used for string matching.

```
SELECT S.age, age1=S.age-5, 2*S.age AS age2 FROM Sailors S WHERE S.sname LIKE 'B_%B'
```

'_' stands for any one character
'%' stands for 0 or more arbitrary characters

More Operations

SQL queries produce new tables

If the results of two queries are set-compatible (same # and types columns) then we can apply logical operations

UNION

INTERSECTION

SET DIFFERENCE (called EXCEPT or MINUS)

Find sids of sailors who have reserved a red or a green boat

UNION: Can be used to compute the union of any two union-compatible sets of tuples (which are themselves the result of SQL queries)

```
SELECT R.sid

FROM Boats B,Reserves R

WHERE R.bid=B.bid AND

(B.color='red' OR B.color='green')
```

VS.

```
SELECT R.sid
FROM Boats B, Reserves R
WHERE R.bid=B.bid AND B.color='red'
UNION SELECT R.sid
FROM Boats B, Reserves R
WHERE R.bid=B.bid AND
B.color='green'
```

Find sids of sailors who have reserved a red and a green boat

If we simply replace OR by AND in the previous query, we get the wrong answer. (Why?)
Instead, could use a self-join:

```
SELECT R1.sid
FROM Boats B1, Reserves R1,
Boats B2, Reserves R2
WHERE R1.sid=R2.sid
AND R1.bid=B1.bid
AND R2.bid=B2.bid
AND (B1.color='red' AND B2.color='green')
```

AND Continued...

Key field!

INTERSECT: discussed in the book. Can be used to compute the intersection of any two *union-compatible* sets of tuples

Also in text: EXCEPT (sometimes called MINUS)
Included in the SQL/92 standard, but some systems do not support them

SELECT S.sid FROM Sailors S, Boats B, Reserves R WHERE S.sid=R.sid AND R.bid=B.bid AND B.color='red' INTERSECT SELECT S.sid FROM Sailors S, Boats B, Reserves R WHERE S.sid=R.sid AND R.bid=B.bid AND B.color='green'

Your turn ...

- Find (the names of) all sailors who are over
 years old
- Find (the names of) all boats that have been reserved at least once
- 3. Find all sailors who have <u>not</u> reserved a red boat (hint: use "EXCEPT")
- 4. Find all pairs of same-color boats
- 5. Find all pairs of sailors in which the <u>older</u> sailor has a <u>lower</u> rating

Find (the names of) all sailors who are over
 years old

```
SELECT S.sname
FROM Sailors S
WHERE S.age > 50
```

2. Find (the names of) all boats that have been reserved at least once

```
SELECT DISTINCT B.bname
FROM Boats B, Reserves R
WHERE R.bid=B.bid
```

3. Find all sailors who have <u>not</u> reserved a red boat

```
SELECT S.sid
FROM Sailors S
EXCEPT
SELECT R.sid
FROM Boats B,Reserves R
WHERE R.bid=B.bid
AND B.color='red'
```

4. Find all pairs of same-color boats



```
SELECT B1.bname, B2.bname
FROM Boats B1, Boats B2
WHERE B1.color = B2.color
AND B1.bid < B2.bid
```

5. Find all pairs of sailors in which the <u>older</u> sailor has a <u>lower</u> rating

```
SELECT S1.sname, S2.sname

FROM Sailors S1, Sailors S2

WHERE S1.age > S2.age

AND S1.rating < S2.rating
```

Nested Queries

powerful feature of SQL:

WHERE clause can itself contain an SQL query!

Actually, so can FROM and HAVING clauses.

Names of sailors who have reserved boat #103

```
SELECT S.sname
FROM Sailors S
WHERE S.sid IN (SELECT R.sid
FROM Reserves R
WHERE R.bid=103)
```

Nested Queries

to find sailors who have not reserved #103, use NOTIN.

To understand semantics of nested queries:

think of a *nested loops* evaluation

for each Sailors tuple

check the qualification by computing the subquery

Nested Queries with Correlation

Find names of sailors who have reserved boat #103

```
SELECT S.sname

FROM Sailors S
WHERE EXISTS (SELECT *

FROM Reserves R
WHERE R.bid=103 AND S.sid=R.sid)
```

EXISTS is another set operator, like IN (also NOT EXISTS)

If EXISTS UNIQUE is used, and * is replaced by *R.bid*, finds sailors with at most one reservation for boat #103.

UNIQUE checks for duplicate tuples in a subquery;

Subquery must be recomputed for each Sailors tuple.

Think of subquery as a function call that runs a query!

More on Set-Comparison Operators

We've already seen IN, EXISTS and UNIQUE. Can also use NOT IN, NOT EXISTS and NOT UNIQUE.

Also available: op ANY, op ALL

Find sailors whose rating is greater than that of some sailor called Horatio:

```
SELECT *
FROM Sailors S
WHERE S.rating > ANY (SELECT S2.rating
FROM Sailors S2
WHERE S2.sname='Horatio')
```

Rewriting INTERSECT Queries Using IN

Find sids of sailors who have reserved both a <u>red and a green</u> boat

```
SELECT R.sid

FROM Boats B, Reserves R

WHERE R.bid=B.bid

AND B.color='red'

AND R.sid IN (SELECT R2.sid

FROM Boats B2, Reserves R2

WHERE R2.bid=B2.bid

AND B2.color='green')
```

Similarly, EXCEPT queries can be re-written using NOT IN. How would you change this to find *names* (not *sids*) of Sailors who've reserved both red and green boats?

Query #3 revisited ...

 Find all sailors who have <u>not</u> reserved a red boat (this time, without using "EXCEPT")

Answer ...

Find all sailors who have <u>not</u> reserved a red boat

Another correct answer ...

Find all sailors who have <u>not</u> reserved a red boat

```
SELECT S.sid

FROM Sailors S

WHERE NOT EXISTS

(SELECT *

FROM Reserves R, Boats B

WHERE R.sid = S.sid

AND R.bid = B.bid

AND B.color = 'red')
```

Division in SQL

Find sailors who have reserved all boats.

```
SELECT S.sname
FROM Sailors S Sailors S such that ...

WHERE NOT EXISTS (SELECT B.bid there is no boat B without ...
FROM Boats B
```

WHERE NOT EXISTS (SELECT R.bid

EDOM Pacative

a Reserves tuple showing S reserved B

FROM Reserves R
WHERE R.bid=B.bid
AND R.sid=S.sid))

Aggregate Operators

Significant extension of relational algebra.

SELECT COUNT (*)
FROM Sailors S

SELECT AVG (S.age) FROM Sailors S WHERE S.rating=10

SELECT COUNT (DISTINCT S.rating)
FROM Sailors S
WHERE S.sname='Bob'

COUNT (*)
COUNT ([DISTINCT] A)
SUM ([DISTINCT] A)
AVG ([DISTINCT] A)
MAX (A)
MIN (A)

single column

Aggregate Operators

```
COUNT (*)
COUNT ( [DISTINCT] A)
SUM ( [DISTINCT] A)
AVG ( [DISTINCT] A)
MAX (A)
MIN (A)
```

single column

SELECT S.sname
FROM Sailors S
WHERE S.rating= (SELECT MAX(S2.rating)
FROM Sailors S2)

SELECT AVG (DISTINCT S.age) FROM Sailors S WHERE S.rating=10

Find name and age of the oldest sailor(s)

The first query is incorrect!

Third query equivalent to second query

allowed in SQL/92 standard, but not supported in some systems.

SELECT S.sname, MAX (S.age) FROM Sailors S

SELECT S.sname, S.age
FROM Sailors S
WHERE S.age =
(SELECT MAX (S2.age)
FROM Sailors S2)

SELECT S.sname, S.age
FROM Sailors S
WHERE (SELECT MAX (S2.age)
FROM Sailors S2)
= S.age

GROUP BY and HAVING

So far, we've applied aggregate operators to all (qualifying) tuples.

Sometimes, we want to apply them to each of several *groups* of tuples.

Consider: Find the age of the youngest sailor for each rating level.

In general, we don't know how many rating levels exist, and what the rating values for these levels are!

Suppose we know that rating values go from 1 to 10; we can write 10 queries that look like this (!):

For
$$i = 1, 2, ..., 10$$
:

SELECT MIN (S.age)
FROM Sailors S
WHERE S.rating = i

Queries With GROUP BY and HAVING

SELECT [DISTINCT] target-list

FROM relation-list

WHERE qualification

GROUP BY grouping-list

[HAVING group-qualification]

Group rows by columns in grouping-list

Use the HAVING clause to restrict which grouprows are returned in the result set

Conceptual Evaluation

- 1. Cross-product of *relation-list*
- 2. Select only tuples that follow the where clause (qualification)
- 3. Partition rows by the value of attributes in *grouping-list*
- 4. Select only groups that follow the *group-qualification*

Expressions in *group-qualification* must have a <u>single value per group</u>! That is, attributes in *group-qualification* must be arguments of an aggregate op or must also appear in the *grouping-list*.

5. One answer tuple is generated per qualifying group.

Find the age of the youngest sailor with age ≥ 18 , for each rating with at least 2 <u>such</u> sailors

SELECT S.rating, MIN (S.age)

FROM Sailors S

WHERE S.age >= 18

GROUP BY S.rating

HAVING COUNT (*) > 1

sid	sname	rating	age
22	dustin	7	45.0
31	lubber	8	55.5
71	zorba	10	16.0
64	horatio	7	35.0
29	brutus	1	33.0
58	rusty	10	35.0

rating	age
1	33.0
7	45.0
7	35.0
8	55.5
10	35.0

rating	m-age	count
1	33.0	1
7	35.0	2
8	55.0	1
10	35.0	1

rating	
7	35.0

Find sailors who've reserved all boats.

Can you do this using Group By and Having?

```
SELECT S.name
```

FROM Sailors S, Reserves R

WHERE S.sid = R.sid

GROUP BY S.name, S.sid

HAVING COUNT(DISTINCT R.bid) =

(Select COUNT (*) FROM Boats)

Note: must have both sid and name in the GROUP BY clause. Why?

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```
SELECT S.name, S.sid

FROM Sailors S, reserves R

WHERE S.sid = R.sid

GROUP BY S.name, S.sid

HAVING COUNT(DISTINCT R.bid) =

(Select COUNT (*) FROM Boats)
```

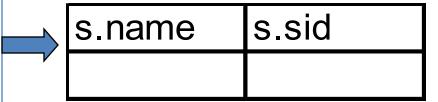
s.name	s.sid	r.sid	r.bid	
Dustin	22	22		101
Lubber	31	22		101
Bob	95	22		101
Dustin	22	95		102
Lubber	31	95		102
Bob	95	95		102

bid	bname	color
101	Interlake	blue
102	Interlake	red
103	Clipper	green
104	Marine	red

Count (*) from boats = 4

s.name	s.sid	bcount
Dustin	22	1
Bob	95	1

Apply having clause to groups



Sorting the Results of a Query

ORDER BY column [ASC | DESC] [, ...]

```
SELECT S.rating, S.sname, S.age
```

FROM Sailors S, Boats B, Reserves R

WHERE S.sid=R.sid AND R.bid=B.bid

AND B.color='red'

ORDER BY S.rating, S.sname;

Sorting the Results of a Query

ORDER BY column [ASC | DESC] [, ...]

```
SELECT S.rating, S.sname, S.age
FROM Sailors S, Boats B, Reserves R
WHERE S.sid=R.sid AND R.bid=B.bid
AND B.color='red'
ORDER BY S.rating, S.sname;
```

Extra reporting power obtained by combining with aggregation.

```
SELECT S.sid, COUNT (*) AS redrescnt
FROM Sailors S, Boats B, Reserves R
WHERE S.sid=R.sid AND R.bid=B.bid
AND B.color='red'
GROUP BY S.sid
ORDER BY redrescnt DESC;
```

Summary: The SQL Query

```
SELECT [DISTINCT] target-list
```

FROM relation-list

WHERE qualification

GROUP BY grouping-list

HAVING group-qualification

ORDER BY attribute-list