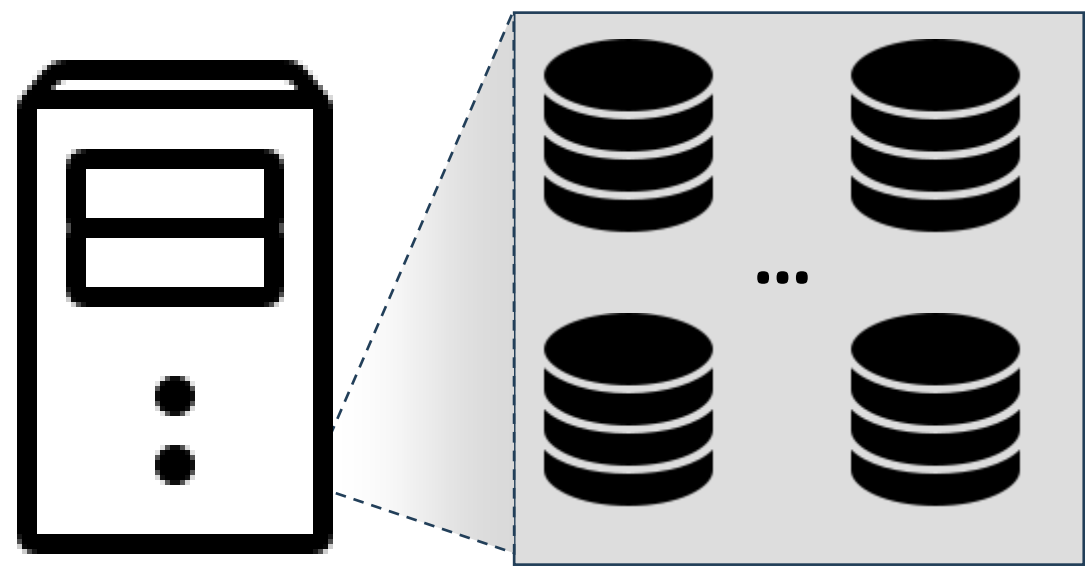


NOCAP: Near-Optimal Correlation-Aware Partitioning Joins

Zichen Zhu Xiao Hu Manos Athanassoulis

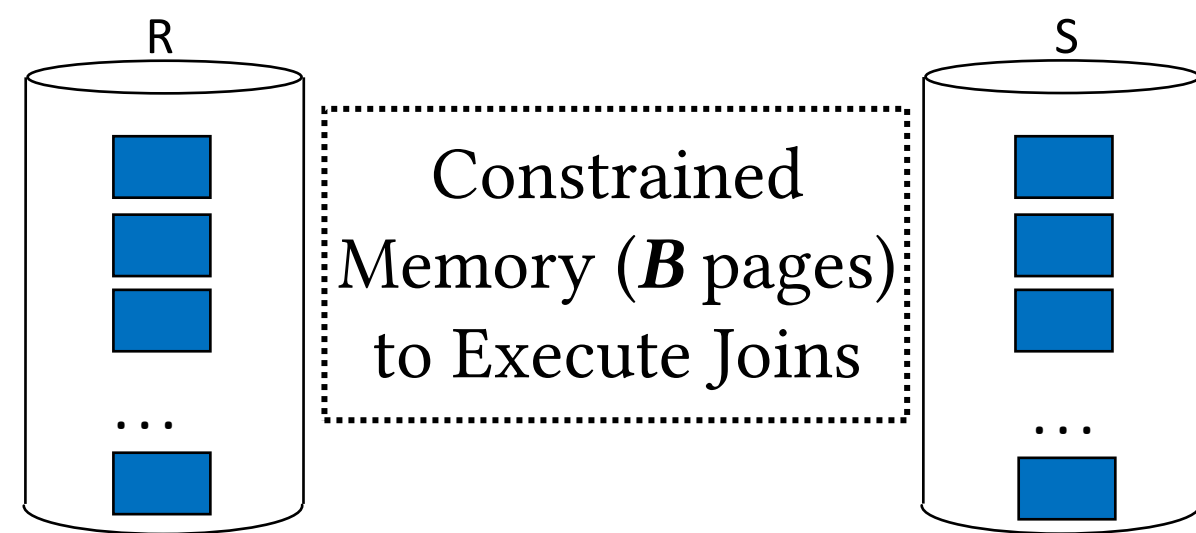
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Storage-based Joins



Multiple DB instances
in the cloud share the
memory resources

Tables may not fit in
memory, which leads to
storage-based joins.

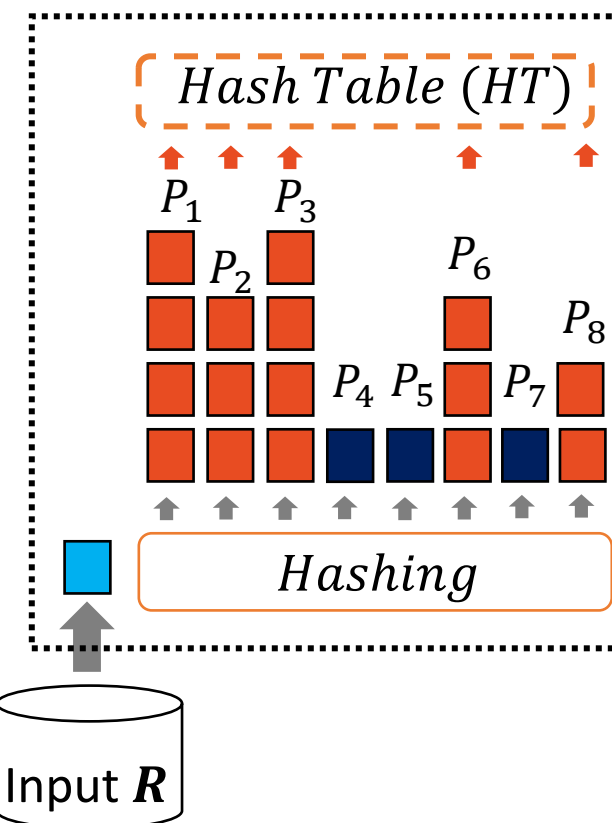


Dynamic Hybrid Hash Join (DHH)

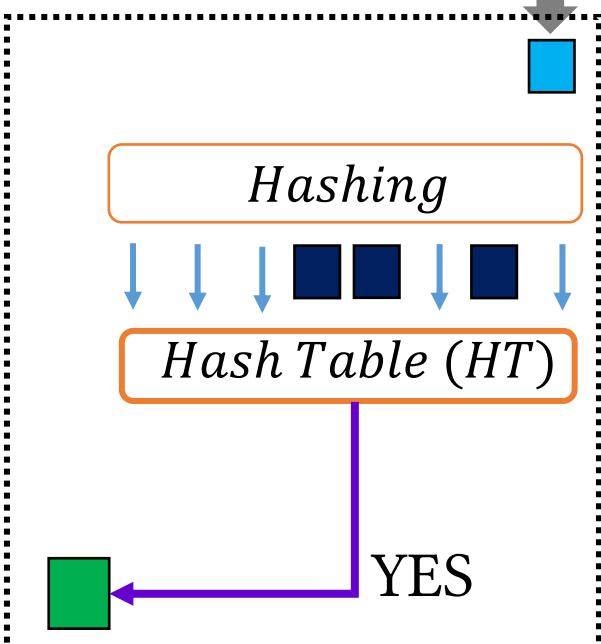
State-of-the-art DB engines use DHH to
implement storage-based joins



Partition R



Partition S



Output Buffer Input Page Pages of in-memory partitions Pages of on-disk partitions

Optimal Correlation-Aware Partitioning (OCAP)

An $n \times (m + 1)$ Boolean matrix P that represents the
partitioning assignment (**list-based partitioning**)

$$P = \begin{bmatrix} 0 & \cdots & 1 \\ \vdots & \ddots & \vdots \\ 1 & \cdots & 0 \end{bmatrix}_{n \times (m+1)}$$

Notation	Meaning
n (n_R)	The number of tuples in relation R
m	The number of partitions on disk
R_1	A partition cached in memory
$Cost(R_j , S_j)$	The read and write cost of R_j, S_j plus the join cost between R_j, S_j

$$\arg \min_{P, m} \sum_{j=2}^{m+1} Cost(||R_j||, ||S_j||)$$

$$\text{s.t. } \forall i \in [n], \sum_{j=1}^{m+1} P_{i,j} = 1$$

$$||R_1|| + m + 2 \leq B$$

$$P_{i,j} \in \{0,1\}, \forall i \in [n], \forall j \in [m + 1]$$

Exponential searching space !

Consecutiveness Monotonicity Divisibility

$O(n^2 \cdot \log B/B)$ for PK-FK joins

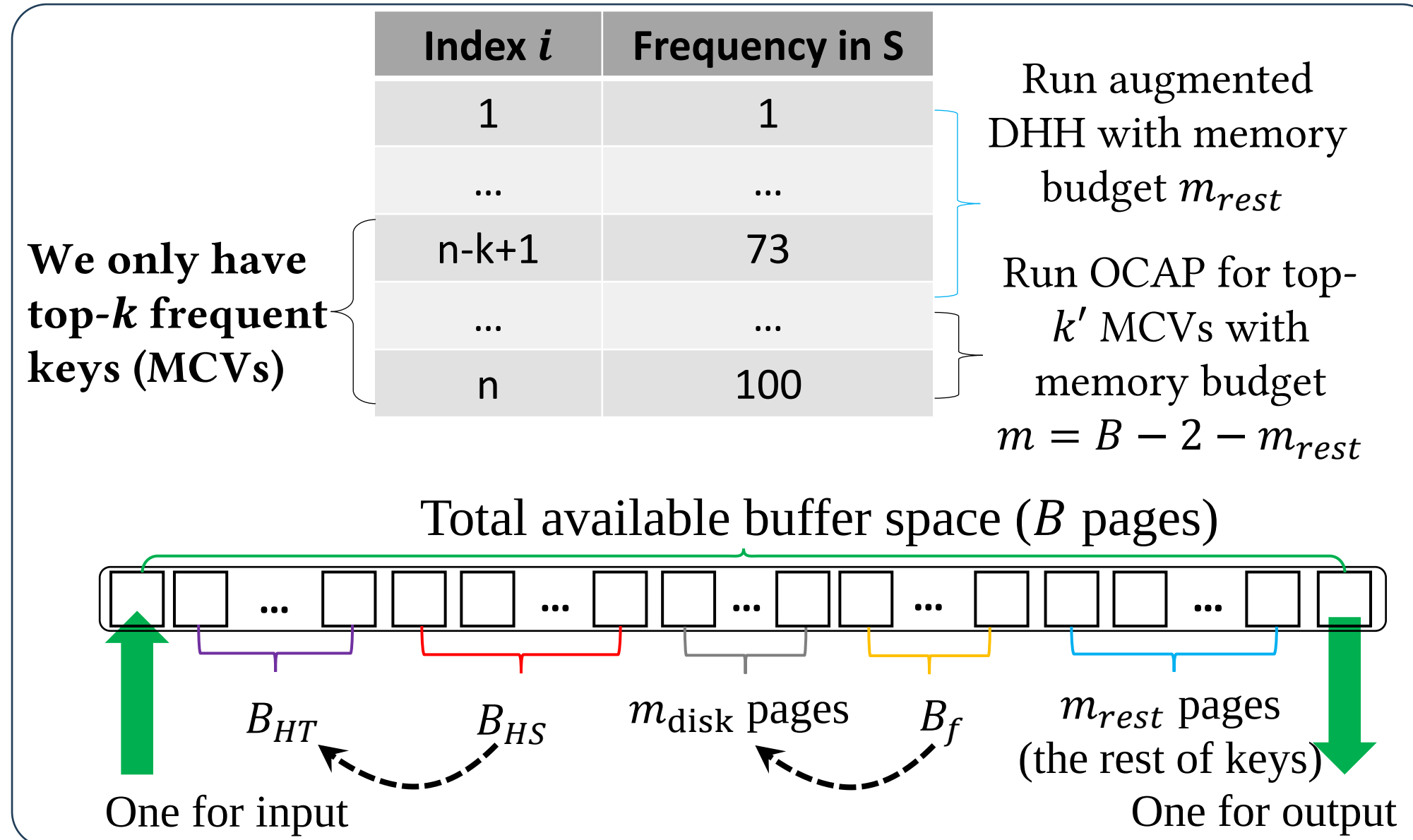
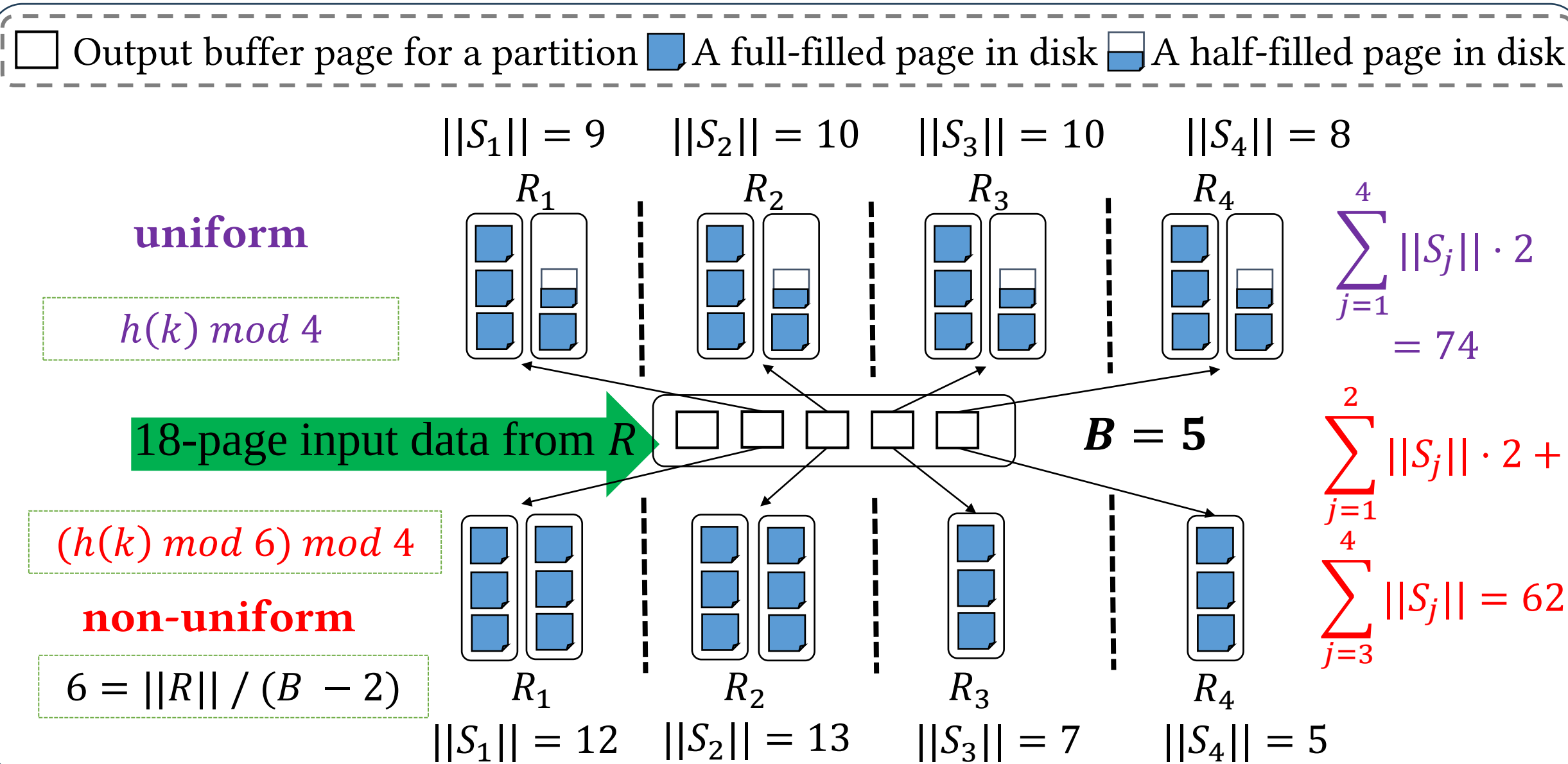


Practical
Challenges

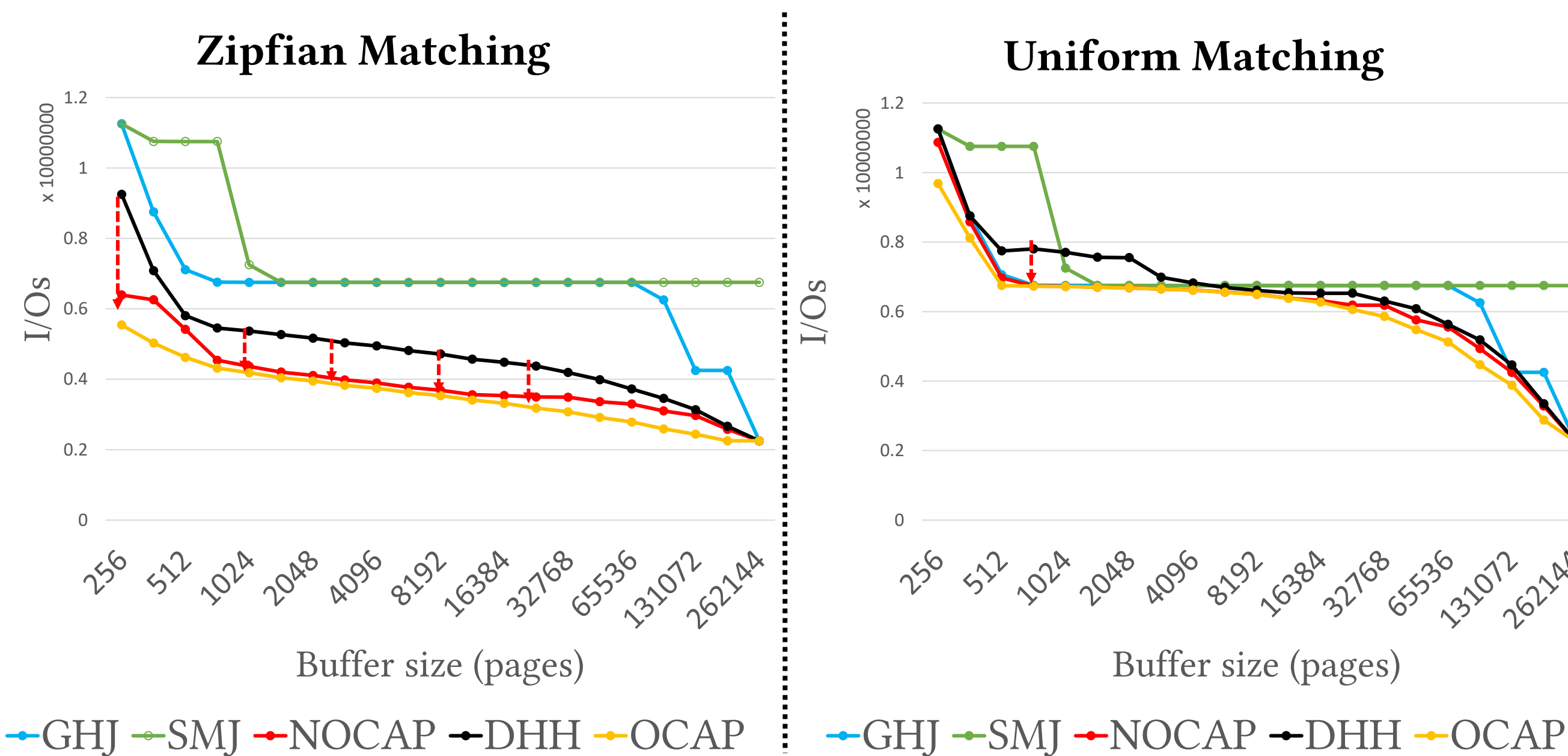
(1) Limited
matching info

(2) List-based
partitioning also
occupies memory

Near-OCAP (NOCAP) = Divisibility + Prioritizing MCVs with Constrained Memory



Experiments

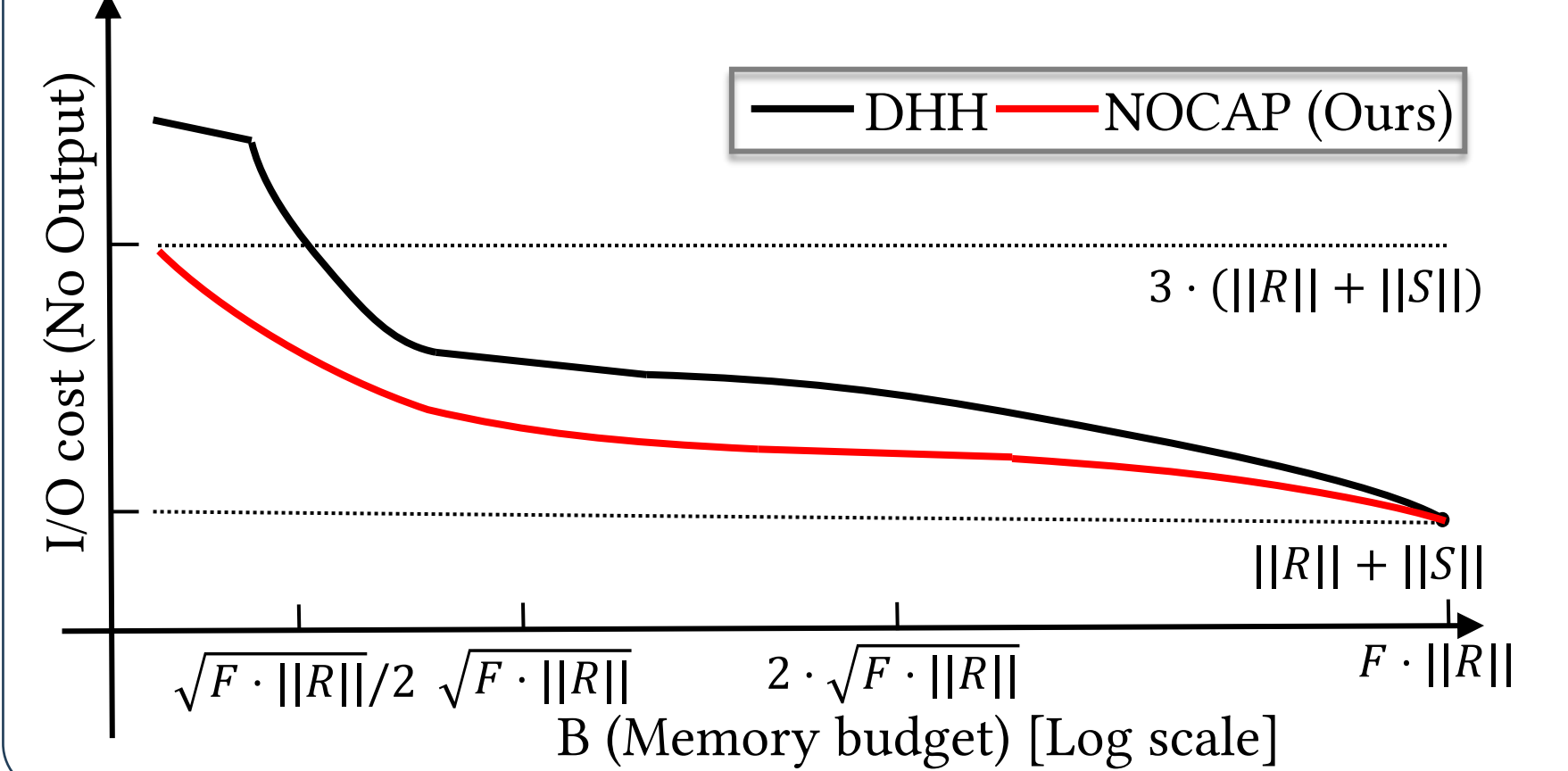


Summary of NOCAP

Up to 30% improvement for skewed join correlation,

Up to 10% even for uniform case!

outperforms DHH under *any* memory budget!



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DISC