Math and Algorithms Review

A firm background in discrete math, algorithms, and mathematical problem-solving / proof-writing will set you up for success in CS 332. You should be comfortable with almost all of the topics listed below and know how to complete the suggested exercises. If one or two of them are rusty, you can review them in the textbooks used for CS 131 (Rosen, Discrete Math and its Applications) and CS 330 (Kleinberg & Tardos, Algorithm Design). However, if many of the topics are unfamiliar, we urge you to take a different course (e.g., 131, 132, 235, 237, 330) to shore up your background.

1 Discrete Math


Exercise 1. Write a short, informal English description of the following sets.

1. \( \{1, 3, 5, 7, \ldots \} \)
2. \( \{\ldots, -4, -2, 0, 2, 4, \ldots \} \)
3. \( \{n : n = 2m \text{ for some } m \in \mathbb{N} \text{ and } n = 3k \text{ for some } k \in \mathbb{N} \} \)
4. \( \{n \in \mathbb{Z} : n = n + 1\} \)

Exercise 2. Write a formal description of each of the following sets.

1. The set containing all integers greater than 5
2. The set containing all natural numbers that are less than 5
3. The set containing the empty set

Exercise 3. Let \( S = \{1, 2\} \) and \( T = \{1, 2, 3\} \). Is \( S \) a subset of \( T \)? Is \( T \) a subset of \( S \)? What are \( S \cap T \) and \( S \cup T \)? What is the power set \( \mathcal{P}(S) \)? What is the Cartesian product \( S \times T \)?

Exercise 4. If \( |A| = a \) (the size of \( A \) is \( a \)) and \( |B| = b \), what is \( |A \times B| \)? What is \( |\mathcal{P}(A)| \)?

Functions and relations. Function (mapping), domain, co-domain, range. Injective (one-to-one), surjective (onto), and bijective mappings. Relation. Equivalence relation (reflexivity, symmetry, transitivity).

Exercise 5. Let \( S = \{1, 2\} \) and \( T = \{1, 2, 3\} \). Give an example of an injective function \( f : S \rightarrow T \). Give an example of a function \( f : S \rightarrow T \) which is not injective. Can a function \( f : S \rightarrow T \) be surjective?

Define the function \( f : S \times T \rightarrow \mathbb{Z} \) by \( f(x, y) = x + y \). What is the domain of \( f \)? What is the range of \( f \)? What is the value \( f(1, 1) \)?

Exercise 6. Define the relation \( \sim \) on the integers by \( a \sim b \) iff \( a - b \) is even. Prove that \( \sim \) is an equivalence relation.

Exercise 7. Consider an undirected graph \( G = (V, E) \) with vertex set \( V = \{1, 2, 3, 4\} \) and edge set \( E = \{\{1, 2\}, \{2, 3\}, \{1, 3\}, \{2, 4\}, \{1, 4\}\} \). Draw the graph \( G \). What is the degree of vertex 1? of vertex 3? Draw a path from vertex 3 to vertex 4. Find a cycle of length 3. Is the graph connected?

Pigeonhole principle. If \( n \) pigeons are placed into \( n - 1 \) pigeonholes, then some pigeonhole must contain more than one pigeon.

Exercise 8. Choose five distinct integers between 0 and 7, inclusive. Show that there exists a pair of them which add up to 7.

Propositional logic. Negation (not), conjunction (and), disjunction (or). Exclusive or (xor), equality, implication. Distributive laws for and/or, De Morgan’s law. Converse, inverse, contrapositive. Quantifiers (\( \exists, \forall \)).

Exercise 9. Negate the following English sentence: Every CS major who is awesome is enrolled in CS 332.

Exercise 10. Show that the following Boolean formulas are equivalent: \( \overline{a} \land (b \lor \overline{c}) \) and \( (\overline{a} \land b) \lor (a \lor c) \).

Exercise 11. Which of the following statements is true? \( \forall x \in \mathbb{N}, \exists y \in \mathbb{N}, y = x + 1 \), or \( \exists x \in \mathbb{N}, \forall y \in \mathbb{N}, y = x + 1 \).

2 Algorithms

Asymptotic notation. Big-Oh \( O \), little-oh \( o \), Big-Omega \( \Omega \), little-omega \( \omega \), Theta \( \Theta \).

Exercise 12. Show that \( 3n^2 + 2n - 1 = O(n^2) \). Show that \( 2^n = o(3^n) \).

Graph algorithms. Breadth-first search, depth-first search, finding connected components.

Dynamic programming. Top-down and bottom-up recursion. DP algorithms for computing Fibonacci numbers, interval scheduling, and Knapsack.

Exercise 13. Design an algorithm for the following problem. Given \( n \) denominations of coins \( \{c_1, \ldots, c_n\} \) and a target value \( V \), find the fewest number of coins needed to make change for \( V \). (Or report that making change for \( V \) is impossible.) Your algorithm should run in time \( O(nV) \).

NP-completeness. Problems in NP, reductions, definition of NP-completeness. Classic NP-complete problems: SAT, clique, 3-coloring, Hamiltonian path. (It’s ok if this is unfamiliar since building the formal theory of NP-completeness is a major topic of 332. However, it will be useful to come into that part of the course with some intuition.)
3 Proof Techniques

Constructions and counterexamples. To prove the statement that some kind of object exists, a constructive proof exhibits such an object. To disprove a statement about all members of some set of objects, it suffices to find a counterexample.

Exercise 14. Prove or give a counterexample to each of the following statements.

1. Every real number is an even integer.
2. For every positive even number $n$, there exists a number $k \geq n$ where $k$ is divisible by 3.

Proof by contraposition. Proving a statement of the form $P \implies Q$ is equivalent to proving that $\neg Q \implies \neg P$.

Exercise 15. Give a contrapositive proof of the following statement. If $n$ is an integer for which $n^2$ is odd, then $n$ is odd.

Proof by contradiction. To prove that a theorem is true, assume instead that the theorem is false and use it to derive a false consequence.

Exercise 16. Show that if $a$ is rational and $ab$ is irrational, then $b$ is irrational.

Proof by induction. Ingredients of an inductive proof: Basis step, induction hypothesis, induction step. Strong induction. Structural induction (i.e., induction to reason about recursively defined objects, e.g., binary trees).

Exercise 17. Prove by induction on $n$ that $2^0 + 2^1 + \cdots + 2^{n-1} + 2^n = 2^{n+1} - 1$ for every $n \in \mathbb{N}$.

Exercise 18. Find the error in the following “proof” that all BU students are from the same city:

Claim: For every $n \geq 1$, every set of $n$ BU students is from the same city.

Proof: By induction on $n$.

Basis step: Let $n = 1$. In any set of one BU student, all students in that set are from the same city.

Induction step: For $n \geq 1$, assume the claim is true for $n$ students and prove it is true for $n + 1$ students. Let $B$ be any set of $n + 1$ BU students. We show that all students in $B$ are from the same city. Remove one student from this set to obtain the set $B_1$. By the induction hypothesis, all the students in $H_1$ are from the same city. Now replace the removed student and remove a different student to obtain the set $H_2$. Again by the induction hypothesis, all the students in $H_2$ are from the same city. Therefore all the students in $H$ are from the same city.