

Lecture Notes 15:**Random Walks****Reading.**

- Spielman Chapter 10

A random walk is a process that begins at a vertex in a graph, and in each time step, moves to an adjacent vertex. Last time, we saw that the second eigenvalue of the normalized Laplacian governs the behavior of one step of a random walk. Today we will see how it governs a random walk's convergence rate, or "mixing time."

1 Random Walks

Let $G = (V, E)$ be an undirected graph. We will primarily be interested in the expected behavior of a random walk, i.e., the distribution of a walk after a certain number of time steps.

Denote a probability distribution on V by its probability mass function $p : V \rightarrow [0, 1]$. Let p_t denote the distribution of a random walk at time step t . We will usually think of starting the walk deterministically from a fixed vertex, so $p_i = \mathbf{1}_{\{v_0\}}$ for some $v_0 \in V$. To sample from the distribution of random walk at time step $t + 1$:

1. Sample a vertex $u \sim p_t$
2. Sample a random vertex $v \sim u$.

Algebraically, this update rule is given by:

$$p_{t+1}(v) = \sum_{(u,v) \in E} \frac{1}{\deg(u)} \cdot p_t(u).$$

It's a useful exercise to check that

$$p_{t+1} = W p_t = W^t p_0$$

where $W = AD^{-1}$ is the *walk matrix* of the graph G .

The walk matrix has a special (right) eigenvector with positive entries,

$$\pi(u) = \frac{\deg(u)}{\sum_{v \in V} \deg(v)} = \frac{\deg(u)}{\|d^{1/2}\|}$$

with eigenvalue 1:

$$W\pi(v) = \sum_{(u,v) \in E} \frac{1}{\deg(u)} \cdot \frac{\deg(u)}{\|d^{1/2}\|} = \frac{\deg(v)}{\|d^{1/2}\|}.$$

This is the *stationary* or *stable* distribution, as it is the distribution on vertices that remains fixed after taking one step of the random walk.

We'd like to establish conditions under which a random walk converges to its stationary distribution, i.e., $p_t \rightarrow \pi$ as $t \rightarrow \infty$. For intuition, suppose ϕ_1, \dots, ϕ_n is a (not necessarily orthonormal) basis of eigenvectors for W with corresponding eigenvalues $1 = w_1 \geq w_2 \geq \dots \geq w_n \geq -1$. Then if $p_0 = \alpha_1 \phi_1 + \dots + \alpha_n \phi_n$, we have

$$W^t p_0 = \alpha_1 w_1^t \phi_1 + \dots + \alpha_n w_n^t \phi_n.$$

As long as $w_1 > w_2$ and $w_n > -1$, this sum will be dominated by the first term $\alpha_1 \phi_1$, with all other terms vanishing exponentially fast in t .

However, things can go (seriously) wrong when either of these eigenvalue conditions fail. Fortunately, these eigenvalue conditions exactly correspond to conditions under which we expect convergence to fail.

Theorem 1 (Consequence of Perron-Frobenius). *Let W be the walk matrix of a graph G with eigenvalues $1 = w_1 \geq w_2 \geq \dots \geq w_n \geq -1$. Then*

- $w_2 = 1$ iff G is disconnected.
- If G is connected, then $w_n = -1$ iff G is bipartite.

I won't prove this, but to understand why it's true, you can read Chapter 4.5 in Spielman. Note that the Perron-Frobenius theorem doesn't directly apply to W since it's not symmetric, so one instead uses the fact that $W = D^{1/2}(D^{-1/2}AD^{-1/2})D^{-1/2}$ is similar to the symmetric normalized adjacency matrix $D^{-1/2}AD^{-1/2}$.

If G is disconnected, then a random walk starting in one component can never reach a different component. If G is bipartite, then a walk will always oscillate between the two halves of the bipartition, so the distribution at time t will depend on the parity of t . Perhaps miraculously, the Perron-Frobenius theorem tells us that these are the only two possible failure modes.

2 Lazy Random Walks

Having to treat bipartite graphs differently is rather annoying, so instead of studying simple random walks, one often studies *lazy random walks*, where the transition operator is given by

$$\widetilde{W} = \frac{1}{2}I + \frac{1}{2}W = \frac{1}{2}I + \frac{1}{2}AD^{-1}.$$

That is, one step of a lazy random walk stays at the same vertex with probability $1/2$ and transitions to a random neighbor with probability $1/2$.

Claim 2. *The eigenvalues of \widetilde{W} are $1 = \omega_1 \geq \omega_2 \geq \dots \geq \omega_n \geq 0$ where $\omega_i = 1 - \nu_i/2$.*

Here, we recall that $0 = \nu_1 \leq \nu_2 \leq \dots \leq \nu_n \leq 2$ are the eigenvalues of the normalized Laplacian $N = I - D^{-1/2}AD^{-1/2}$.

Proof. Let f be an eigenvector of N with eigenvalue ν . We will show that $D^{1/2}f$ is an eigenvector of \widetilde{W} with eigenvalue $1 - \nu/2$. To see this, we calculate:

$$\begin{aligned}\widetilde{W}D^{1/2}f &= \frac{1}{2}D^{1/2}f + \frac{1}{2}AD^{-1/2}f \\ &= \frac{1}{2}D^{1/2}f + \frac{1}{2}D^{1/2}(I - N)f \\ &= D^{1/2}f - \frac{1}{2}\nu D^{1/2}f.\end{aligned}$$

□

Now the heuristic calculation from before goes through, using the ω 's in place of the w 's, as long as $\omega_2 < 1 \iff \nu_2 > 0 \iff G$ is connected.

3 Convergence Rate

Theorem 3. For any start vertex s , letting $p_0 = \mathbf{1}_{\{s\}}$, and any vertex v and $t \geq 1$, we have

$$|p_t(v) - \pi(v)| \leq \sqrt{\frac{\deg(v)}{\deg(s)}} \cdot \omega_2^t.$$

Proof. Let $p_0 = \alpha_1 D^{1/2}f_1 + \dots + \alpha_n D^{1/2}f_n$, where $D^{1/2}f_1, \dots, D^{1/2}f_n$ are the eigenvectors for \widetilde{W} corresponding to eigenvalues $1 = \omega_1 \geq \omega_2 \geq \dots \geq \omega_n \geq 0$. Then

$$p_t = \alpha_1 D^{1/2}f_1 + \sum_{i=2}^n \omega_i^t \alpha_i D^{1/2}f_i = \pi + \sum_{i=2}^n \omega_i^t \alpha_i D^{1/2}f_i.$$

Write

$$\begin{aligned}p_t(v) &= \langle \mathbf{1}_{\{v\}}, p_t \rangle \\ &= \langle \mathbf{1}_{\{v\}}, \pi + \sum_{i=2}^n \omega_i^t \alpha_i D^{1/2}f_i \rangle \\ &= \pi(v) + \sqrt{\deg(v)} \sum_{i=2}^n \omega_i^t \alpha_i f_i(v),\end{aligned}$$

so our goal is to bound the magnitude of the term on the right.

To do this, we observe that the coefficients in our basis decomposition of p_0 satisfy

$$\alpha_i = \langle f_i, D^{-1/2}\mathbf{1}_{\{s\}} \rangle = (\deg(s))^{-1/2} f_i(s)$$

since the f_i 's are orthonormal. Plugging this in gives

$$\sqrt{\deg(v)} \sum_{i=2}^n \omega_i^t \alpha_i f_i(v) = \sqrt{\frac{\deg(v)}{\deg(s)}} \sum_{i=2}^n \omega_i^t f_i(s) f_i(v).$$

The sum on the right is bounded from above by

$$\omega_2^t \sum_{i=2}^n |f_i(s)| \cdot |f_i(v)|.$$

By Cauchy-Schwarz, this is at most

$$\omega_2^t \sqrt{\sum_{i=1}^n f_i(s)^2} \sqrt{\sum_{i=1}^n f_i(v)^2} \leq \omega_2^t$$

since the orthogonal matrix with columns f_1, \dots, f_n also has orthonormal rows. □