



Randomness in Computing

CS
537

LECTURE 1

Randomness in Computing

- Course information
- Verifying polynomial identities
- Probability Amplification
- Probability Review

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- 1. Course staff**
- 2. Course website(s)**
- 3. Piazza bonus**
- 4. Prerequisites**
- 5. Textbook(s)**
- 6. Syllabus**
- 7. Homework logistics**
- 8. Collaboration policy**
- 9. Exams and grading**

Tips for the course

- Concepts in this course take some time to sink in: be careful not to fall behind.
- Do the assigned reading on each topic before the corresponding lecture.
- Attend the lectures: some material will be presented on the blackboard (and some of it is not in the book).
- Attend the discussions: practice problem solving.
- Take advantage of office hours.
- Be active in lectures and on piazza.
- Allocate lots of time for the course: comparable to a project course, but spread more evenly.

Tips for the course: HW

- Start working on HW early.
- Spread your HW time over multiple days.
- You can work in groups (up to 4 people), but spend 1-2 hours thinking about it on your own before your group meeting.

Tips: learning problem solving

To learn problem solving, you have to do it:

- Try to think how you would solve any presented problem before you read/hear the answer.
- Do exercises in addition to HW.

Tips: how to read a math text

- Not like reading a mystery novel.
- The goal is not to get the answers, but to learn the techniques.
- Always try to foresee what is coming next.
- Always think how you would approach a problem before reading the solution.
- This applies to things that are not explicitly labeled as problems.

Skills we will work on

- Mathematical reasoning
- Expressing your ideas
 - abstractly (suppress inessential details)
 - precisely (rigorously)
- Probabilistic thinking
- Algorithmic thinking
- Problem solving
- Having **FUN** with all of the above!!!



- To speed up algorithms.
- To enable new applications:
 - Symmetry breaking in distributed algorithms, cryptography, privacy, online games and gambling.
- To simulate real world events in physical systems: model them as happening randomly.
- To analyze algorithms when data is generated from some distribution:
 - learning theory, data compression.
- To analyze algorithms when errors happen randomly
 - error-correcting codes.
- Analyzing statistics from sampling.

- $(x + 1)(x - 2)(x + 3)(x - 4)(x + 5)(x - 6) \equiv? x^6 - 7x + 37$

Task: Given two polynomials $F(x)$ and $G(x)$, verify if $F(x) \equiv G(x)$.

Idea 1 (deterministic): Convert both polynomials to canonical form $\sum_{i=0}^d c_i x^i$.

- If $F(x)$ is given as $\prod_{i=1}^d (x - a_i)$,
conversion by consecutively multiplying monomials requires $\Theta(d^2)$ multiplications of coefficients.

Faster with Fourier Transform

Task: Given two polynomials $F(x)$ and $G(x)$, verify if $F(x) \equiv G(x)$.

Idea 2 (randomized): Evaluate the polynomials on random integers.

Let $d = \max$ degree of $F(x)$ and $G(x)$

1. Pick r uniformly from $\{1, \dots, 100d\}$.
2. Compute $F(r)$ and $G(r)$.
3. **reject** if $F(r) \neq G(r)$; o. w. **accept**.

$O(d)$ ops for product form

Error Analysis: Probability of accepting incorrectly

1-sided error

Fundamental Theorem of Algebra

A polynomial of degree d has at most d roots.

$\Pr[\text{error}] =$

Probability space has three components

- Sample space Ω
- Family of allowable events $E \subseteq \Omega$
- A probability function \Pr that maps events E to \mathbb{R} such that
 - $\Pr(E) \in [0,1]$ for any event E ;
 - $\Pr(\Omega) = 1$;
 - For any finite or countable sequence of pairwise disjoint events E_1, E_2, \dots ,

$$\Pr\left(\bigcup_{i \geq 1} E_i\right) = \sum_{i \geq 1} \Pr(E_i).$$

Review: Inclusion-Exclusion Principle

For any two events E_1 and E_2 ,

$$\Pr(E_1 \cup E_2) = \Pr(E_1) + \Pr(E_2) - \Pr(E_1 \cap E_2).$$

For any three events E_1, E_2 and E_3 ,

$$\begin{aligned} \Pr(E_1 \cup E_2 \cup E_3) = & \Pr(E_1) + \Pr(E_2) + \Pr(E_3) \\ & - \Pr(E_1 \cap E_2) - \Pr(E_1 \cap E_3) - \Pr(E_2 \cap E_3) \\ & + \Pr(E_1 \cap E_2 \cap E_3) \end{aligned}$$

For any n events E_1, E_2, \dots, E_n ,

$$\begin{aligned} \Pr\left(\bigcup_{i \in [n]} E_i\right) = & \sum_{i \in [n]} \Pr(E_i) - \sum_{i < j} \Pr(E_i \cap E_j) \\ & + \sum_{i < j < k} \Pr(E_i \cap E_j \cap E_k) - \dots + (-1)^{n+1} \Pr\left(\bigcap_{i \in [n]} E_i\right). \end{aligned}$$

Union Bound

For finite or countable sequence of events E_1, E_2, \dots ,

$$\Pr\left(\bigcup_{i \geq 1} E_i\right) \leq \sum_{i \geq 1} \Pr(E_i).$$

Probability Amplification

- Our algorithm for verifying polynomial identities accepts incorrectly with probability $\leq \frac{1}{100}$

Idea: Repeat the algorithm and accept if all iterations accept.

$$\begin{aligned} & \Pr[\text{error in all } k \text{ iterations}] \\ & \leq \left(\frac{1}{100}\right)^k \end{aligned}$$

Independent events

- Two events E_1 and E_2 are **independent** if

$$\Pr(E_1 \cap E_2) = \Pr(E_1) \cdot \Pr(E_2).$$

- Events E_1, \dots, E_n are **mutually independent** if,

for every subset $I \in [n]$,

$[n]$ denotes $\{1, 2, \dots, n\}$

$$\Pr\left(\bigcap_{i \in I} E_i\right) = \prod_{i \in I} \Pr(E_i).$$

- If all pairs of events among E_1, \dots, E_n are independent then E_1, \dots, E_n are **pairwise independent**.

- Pairwise independence does not necessarily imply mutual independence!



- We toss a fair coin twice.
 - Let A be the event that the 1st flip is HEADS.
 - Let B be the event that the 2nd flip is HEADS.
 - Let C be the event that both flips are the same.
1. Are events A, B, C pairwise independent?
 2. Are they mutually independent?

Conditional Probability

The **conditional probability** of event E given event F is

$$\Pr(E \mid F) = \frac{\Pr(E \cap F)}{\Pr(F)}.$$

Well defined only if $\Pr(F) \neq 0$

- When E and F are independent,

$$\Pr(E \mid F) = \frac{\Pr(E \cap F)}{\Pr(F)} = \frac{\Pr(E) \cdot \Pr(F)}{\Pr(F)} = \Pr(E).$$



- **Event E**: the numbers on two dice sum to 8.
- **Event F**: the numbers on two dice are both even.

What is the probability of E given that F occurred, $\Pr(E|F)$?

- A. Less than $1/3$
- B. $1/3$
- C. Greater than $1/3$, but less than $2/3$
- D. $2/3$
- E. Greater than $2/3$





Card dealing

We deal two cards. What is the probability that the second card is an ace, given that the first is an ace?

- A. $3/52$
- B. $3/51$
- C. $4/52$
- D. $5/52$
- E. None of the answers above are correct.

For any two events E_1 and E_2 ,

$$\Pr(E_1 \cap E_2) = \Pr(E_1) \cdot \Pr(E_2|E_1).$$

For all events E_1, \dots, E_n ,

$$\Pr(\bigcap_{i=1}^n E_i) = \Pr(E_1) \cdot \Pr(E_2|E_1) \cdot \dots \cdot \Pr(E_n | \bigcap_{i=1}^{n-1} E_i)$$

Sampling without replacement

- Let E_i be the event that we choose a root in iteration i

$$\begin{aligned} & \Pr[\text{error in all } k \text{ iterations}] \\ &= \Pr[E_1 \cap \dots \cap E_k] \\ &= \Pr[E_1] \cdot \Pr[E_2|E_1] \cdot \dots \cdot \Pr[E_k|E_1 \cap \dots \cap E_{k-1}] \end{aligned}$$

- It is 0 if $k > d$.

- If $k \leq d$, then
$$\Pr[E_j|E_1 \cap \dots \cap E_{j-1}] = \frac{d - (j - 1)}{100d - (j - 1)}$$

$$\Pr[\text{error in all } k \text{ iterations}] \leq \left(\frac{1}{100}\right)^k$$

