

Randomness in Computing



LECTURE 13 Last time

- Finished routing on hypercube
- Balls-and-Bins model
 - Birthday Paradox
- Application: Bucket Sort **Today**
- Poisson distribution
- Poisson approximation



m balls into n bins

• The probability that bin 1 is empty is



• The probability
$$p_k$$
 that bin 1 has k balls is

$$p_k = \binom{m}{k} \left(\frac{1}{n}\right)^k \left(1 - \frac{1}{n}\right)^{m-k} = \frac{m!}{k! (m-k)!} \frac{1}{n^k} \left(1 - \frac{1}{n}\right)^{m-k}$$

$$= \frac{1}{k!} \left(\frac{m}{n} \cdot \frac{m-1}{n} \dots \frac{m-k+1}{n}\right) \left(1 - \frac{1}{n}\right)^{m-k}$$



• A Poisson random variable with parameter μ is given by the following distribution on j = 0, 1, 2, ...

Taylor expansion: $e^x = \sum_{j=0}^{\infty} \frac{x^j}{j!}$

$$\Pr[X = j] = \frac{e^{-\mu}\mu^j}{j!}$$

Check that probabilities sum to 1: Tay

$$\sum_{j=0}^{\infty} \Pr[X=j] = \sum_{j=0}^{\infty} \frac{e^{-\mu} \mu^j}{j!} =$$

• The expectation of a Poisson R.V. X is $\mathbb{E}[X] =$

$$var[X] = \mu$$
 (See Ex. 5.5)



Theorem

Let X and Y be independent Poisson RVs with means μ_X and μ_Y .

Then X + Y is a Poisson RV with mean $\mu_X + \mu_Y$.

CS 537 Chernoff Bounds for Poisson RVs

Theorem. Let X be a Poisson RV with mean μ .

• (upper tail, additive) If $x > \mu$, then

$$\Pr[X \ge x] \le \frac{e^{-\mu}(e\mu)^x}{x^x}.$$

- (lower tail, additive) If $x < \mu$, then $\Pr[X \le x] \le \frac{e^{-\mu}(e\mu)^x}{x^x}.$
- (upper tail, multiplicative) For any $\delta > 0$,

$$\Pr[X \ge (1+\delta)\mu] \le \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{\mu}$$

• (lower tail, multiplicative) For any $\delta \in (0,1)$, $\Pr[X \le (1-\delta)\mu] \le \left(\frac{e^{-\delta}}{(1-\delta)^{1-\delta}}\right)^{\mu}.$

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CS Poisson Distribution is Limit of Binomial **537** Distribution

TheoremLet $X_n \sim Bin(n, p)$, where p is a function of n and $\lim_{n \to \infty} np = \mu$,a constant independent of n.Then, for all fixed k, $\lim_{n \to \infty} \Pr[X_n = k] = \frac{e^{-\mu}\mu^k}{k!}$.

• Applies to balls-and-bins if m = nc.

CS 537 The Poisson Approximation

- The Balls-and-Bins model has dependences.
- E.g. if Bin 1 is empty, then Bin 2 is less likely to be empty.
- The Poisson Approximation gets rid of dependencies.

CS 537 Poisson approximation

- The Balls-and-Bins model has dependences.
- The Poisson approximation gets rid of dependencies.

• *m* balls into *n* bins u.i.r.
For
$$i \in [n]$$
, let
(real world) $X_i^{(m)} = \#$ of balls in bin *i*
(Poisson world) $Y_i^{(m)} \sim Poisson(\mu)$, where $\mu = \frac{m}{n}$ and
 $Y_i^{(m)}$ are mutually independent.

• If we condition the Poisson distribution on producing exactly *k* balls, then it's the same as the distribution resulting from throwing *k* balls into *n* bins.

CS Poisson Distribution Conditioned on537 getting k balls



CS Poisson Distribution Conditioned on 537 getting k balls



CS Approximating a function537 of the loads of the bins



- Fact (*Stirling's formula*): $n! \sim \sqrt{2\pi n} \left(\frac{n}{\rho}\right)^n$
- Bounds for all $n \in \mathbb{N}$: $\sqrt{2\pi} \sqrt{n} \left(\frac{n}{e}\right)^n \le n! \le e \sqrt{n} \left(\frac{n}{e}\right)^n$

CS Approximating a function537 of the loads of the bins

$$\begin{array}{l} \hline \mathsf{Poisson Approximation Theorem} \\ \mbox{Let } f(x_1, \dots, x_n) \geq 0 \mbox{ for all } x_1, \dots, x_n \in \{0, 1, 2, \dots\}. \mbox{ Then} \\ & \mathbb{E}\left[f\left(X_1^{(m)}, \dots, X_n^{(m)}\right)\right] \leq e\sqrt{m} \cdot \mathbb{E}\left[f\left(Y_1^{(m)}, \dots, Y_n^{(m)}\right)\right]. \\ \hline \mathsf{Proof: } \mathbb{E}\left[f\left(Y_1^{(m)}, \dots, Y_n^{(m)}\right)\right] \xrightarrow{\mathsf{Poisson case}} \\ \hline \mathsf{Proof: } \mathbb{E}\left[f\left(Y_1^{(m)}, \dots, Y_n^{(m)}\right)\right] \xrightarrow{\mathsf{Poisson case}} \\ \hline \mathsf{Proof: } \mathbb{E}\left[f\left(Y_1^{(m)}, \dots, Y_n^{(m)}\right)\right] \sum_{i \in [n]} Y_i^{(m)} = k\right] \cdot \Pr\left[\sum_{i \in [n]} Y_i^{(m)} = k\right] \\ \hline \mathbb{E}\left[f\left(Y_1^{(m)}, \dots, Y_n^{(m)}\right)\right] \sum_{i \in [n]} Y_i^{(m)} = m\right] \cdot \Pr\left[\sum_{i \in [n]} Y_i^{(m)} = m\right] \\ & = \mathbb{E}\left[f\left(X_1^{(m)}, \dots, X_n^{(m)}\right)\right] \qquad \cdot \Pr[Y = m] \end{array}$$

CS Approximating a function **537** of the loads of the bins

Poisson Approximation Theorem

Let $f(x_1, \dots, x_n) \ge 0$ for all $x_1, \dots, x_n \in \{0, 1, 2, \dots\}$. Then $\mathbb{E}\left[f\left(X_1^{(m)}, \dots, X_n^{(m)}\right)\right] \le e\sqrt{m} \cdot \mathbb{E}\left[f\left(Y_1^{(m)}, \dots, Y_n^{(m)}\right)\right].$

- Poisson case: # of balls in each bin is independent Poisson $\left(\frac{m}{n}\right)$
- Corollary. Any event that has probability *p* in the Poisson case has probability ≤ *p* · *e*√*m* in the exact case.
 Proof: Let *X* be the indicator for that event.
 Then E[X] is the probability that event occurs.
- Improvements to Theorem and Corollary If $\mathbb{E}\left[f\left(X_{1}^{(m)}, \dots, X_{n}^{(m)}\right)\right]$ is monotonically nonincreasing (or nondecreasing) in *m*, then $e\sqrt{m}$ can be changed to 2.



n balls into n bins

sufficiently large

- Before (by Chernoff): $\Pr\left[MaxLoad > \frac{3\ln n}{\ln \ln n}\right] \le \frac{1}{n}$ for s.1. n
- Theorem. $\Pr\left[MaxLoad < \frac{\ln n}{\ln \ln n}\right] \le \frac{1}{n}$ for s.l. nProof: Let $M = \frac{\ln n}{\ln \ln n}$