



TEL AVIV UNIVERSITY

Information Security – Theory vs. Reality

0368-4474, Winter 2015-2016

Lecture 3:

Power analysis, correlation power analysis

Lecturer:

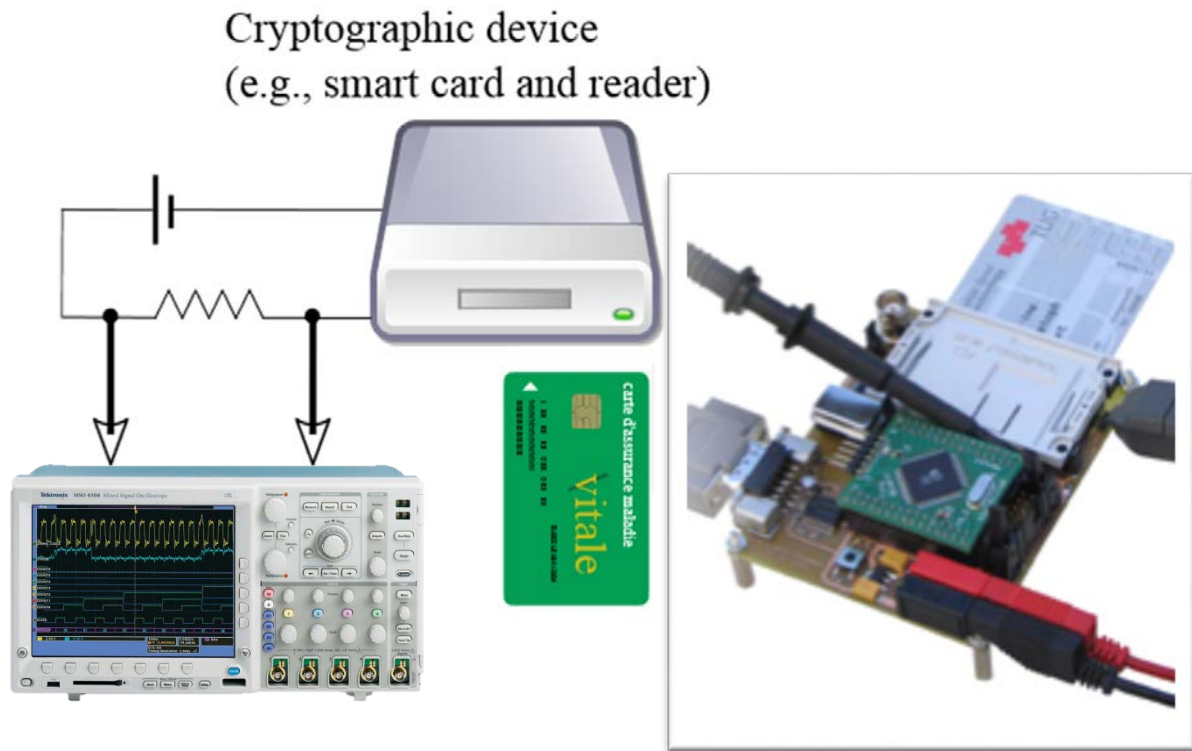
Eran Tromer

Power Analysis

Simple Power Analysis
Correlation Power Analysis
Differential Power Analysis

Power analysis

- Power analysis: measure device's power consumption.



מחפשים מטען?

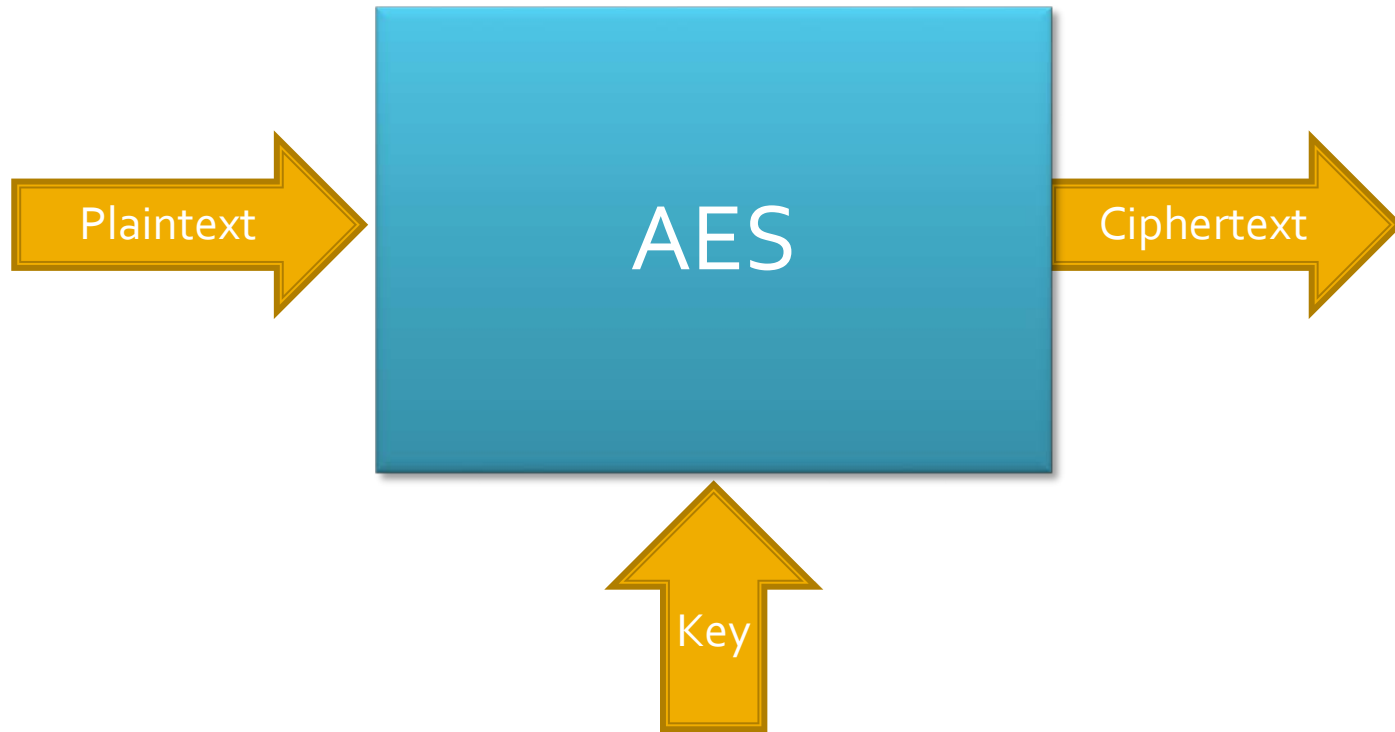
מעכשיו אפשר להשאל במרכז השירות בגילמן
מטענים בחינם.

השירות מתאים לבעלי חיבור Micro USB ואייפון 4 ומטה.
המטענים מתחברים באמצעות מחשב בחיבור usb או בחיבור לנקודת חשמל.
שעות פעילות: 9:00-19:00
טל': 03-6409200
מייל: kabalas@posttau.ac.il
השירות מותנה בהשארת פיקדון של תעודת סטודנט והחזרה עד סוף היום

הסטודנטים



The AES Cipher



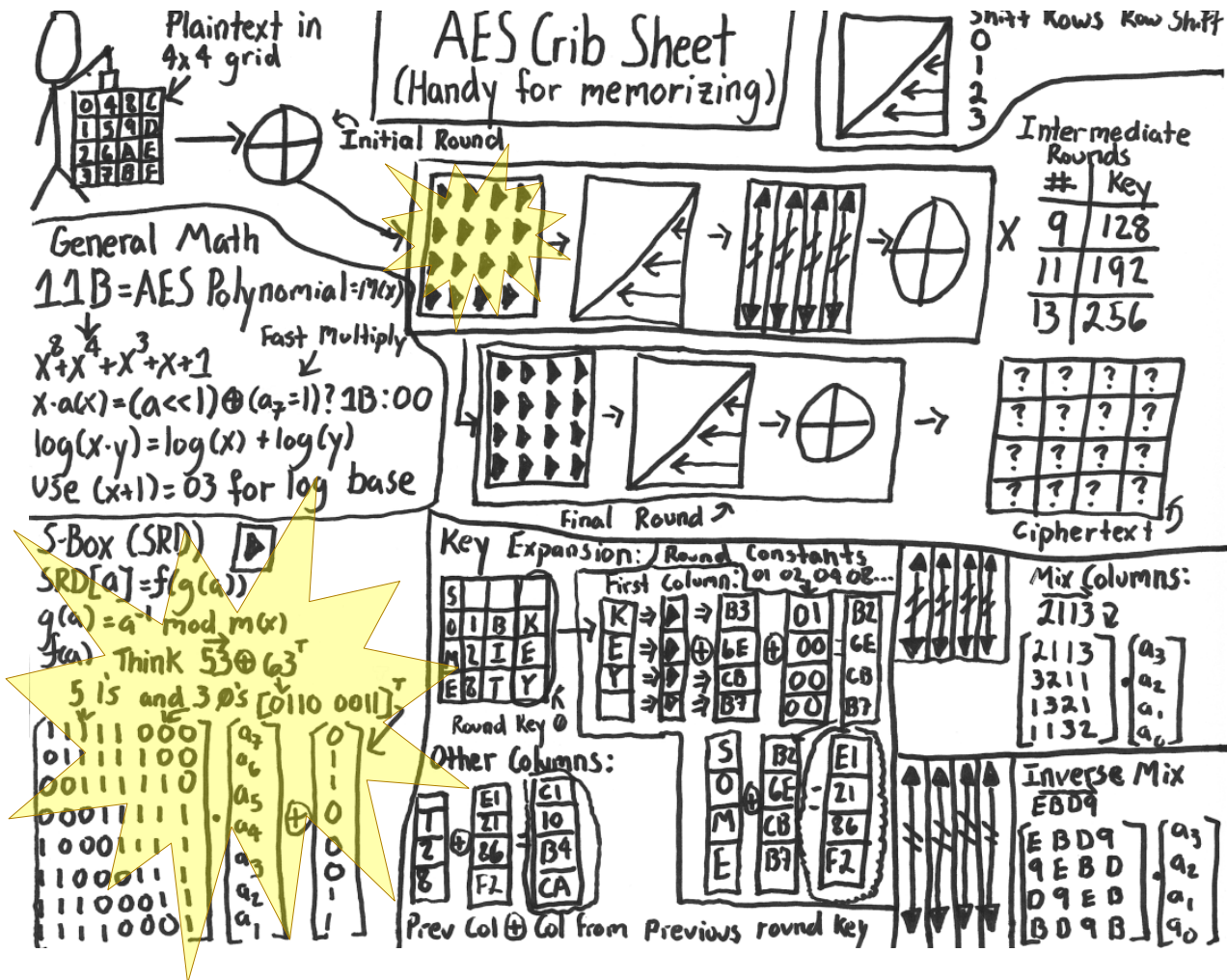
AES symmetric cipher: Lookup-based implementation

```
char p[16], k[16];           // plaintext and key
int32 Col[4];                // intermediate state
const int32 T0[256],T1[256],T2[256],T3[256]; // lookup tables
...

/* Round 1 */
Col[0]← T0[p[ 0]⊕k[ 0]] ⊕ T1[p[ 5]⊕k[ 5]] ⊕
        T2[p[10]⊕k[10]] ⊕ T3[p[15]⊕k[15]];
Col[1]← T0[p[ 4]⊕k[ 4]] ⊕ T1[p[ 9]⊕k[ 9]] ⊕
        T2[p[14]⊕k[14]] ⊕ T3[p[ 3]⊕k[ 3]];
Col[2]← T0[p[ 8]⊕k[ 8]] ⊕ T1[p[13]⊕k[13]] ⊕
        T2[p[ 2]⊕k[ 2]] ⊕ T3[p[ 7]⊕k[ 7]];
Col[3]← T0[p[12]⊕k[12]] ⊕ T1[p[ 1]⊕k[ 1]] ⊕
        T2[p[ 6]⊕k[ 6]] ⊕ T3[p[11]⊕k[11]];
...
```



AES symmetric cipher: "Algebraic" implementation



AES symmetric cipher: “Algebraic” implementation in code

```
void AESEncrypt( u8 input[16], u8 output[16] ) {  
[...]  
    for (r=1; r<=9; r++)  
    {  
        ByteSub(state);  
        ShiftRow(state);  
        MixColumn(state);  
        KeyAdd(state, roundKeys, r);  
    }  
[...]
```

Source: http://users.ece.utexas.edu/~gerstl/ee382v-ics_f09/soc/tutorials/System_C_Code_Examples_2/date04_examples/cosimulate/sw_only/

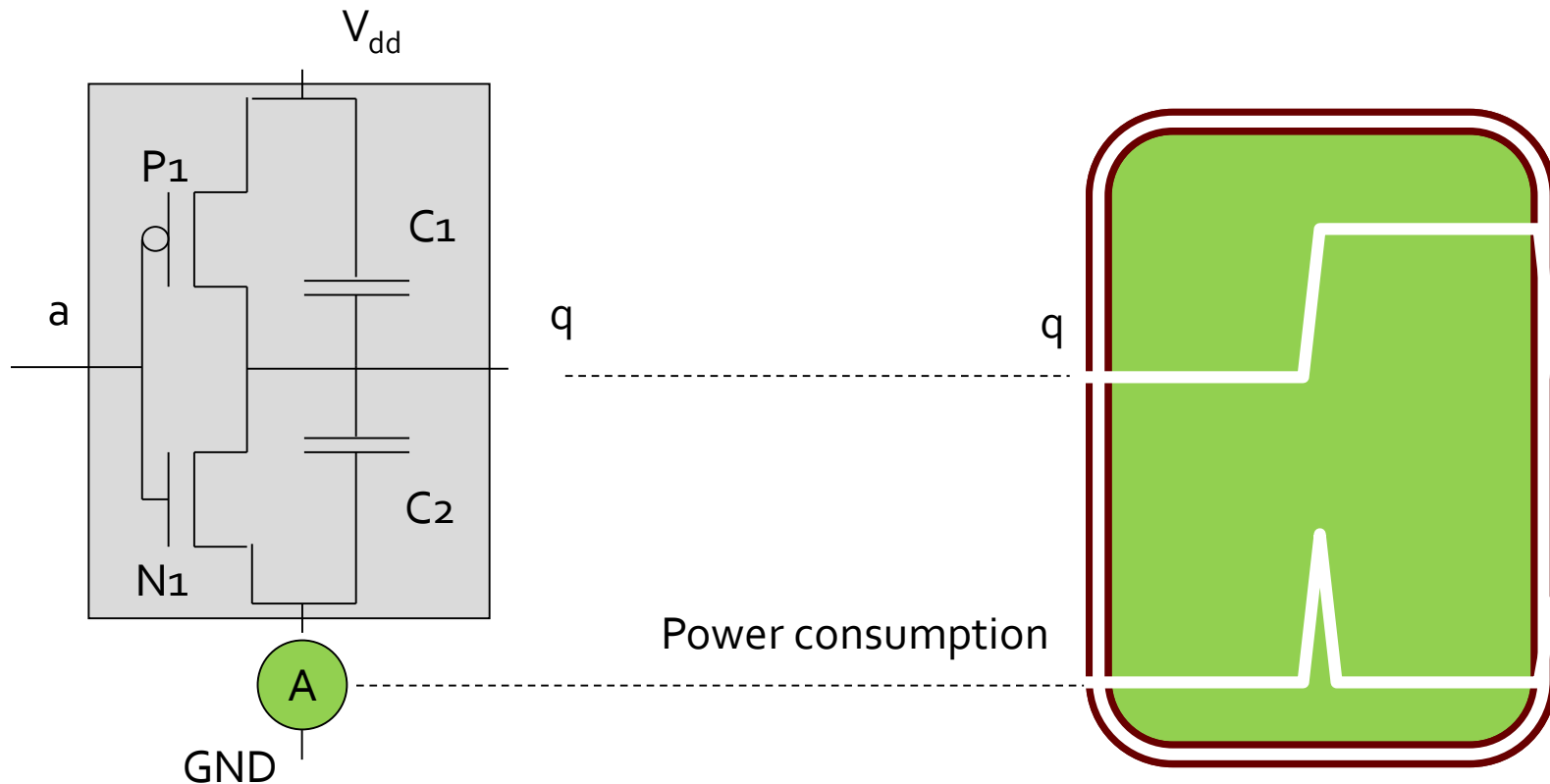


Theory of power analysis

- Power consumption is **variable**
- Power consumption depends on **instruction**
- Power consumption depends on **data**



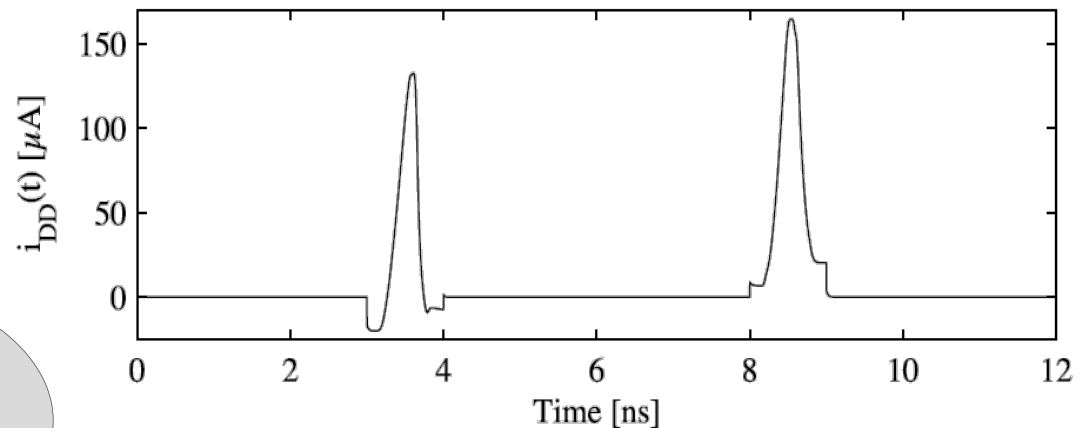
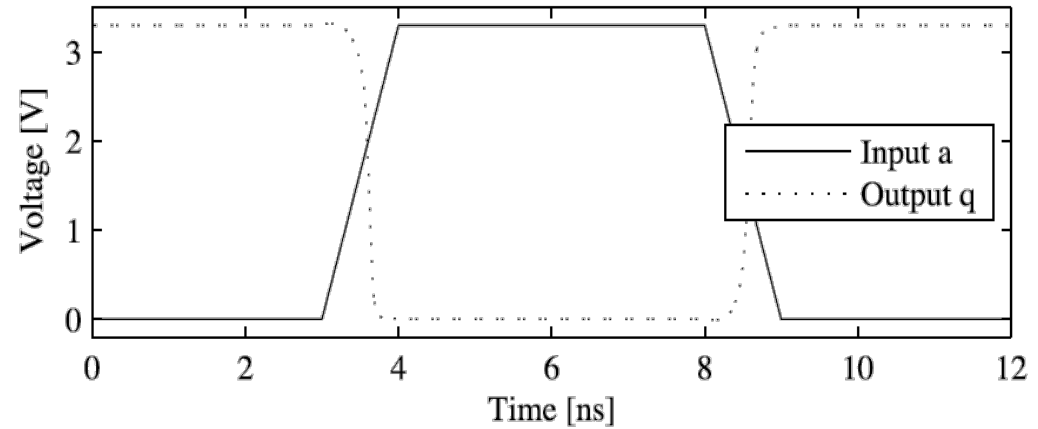
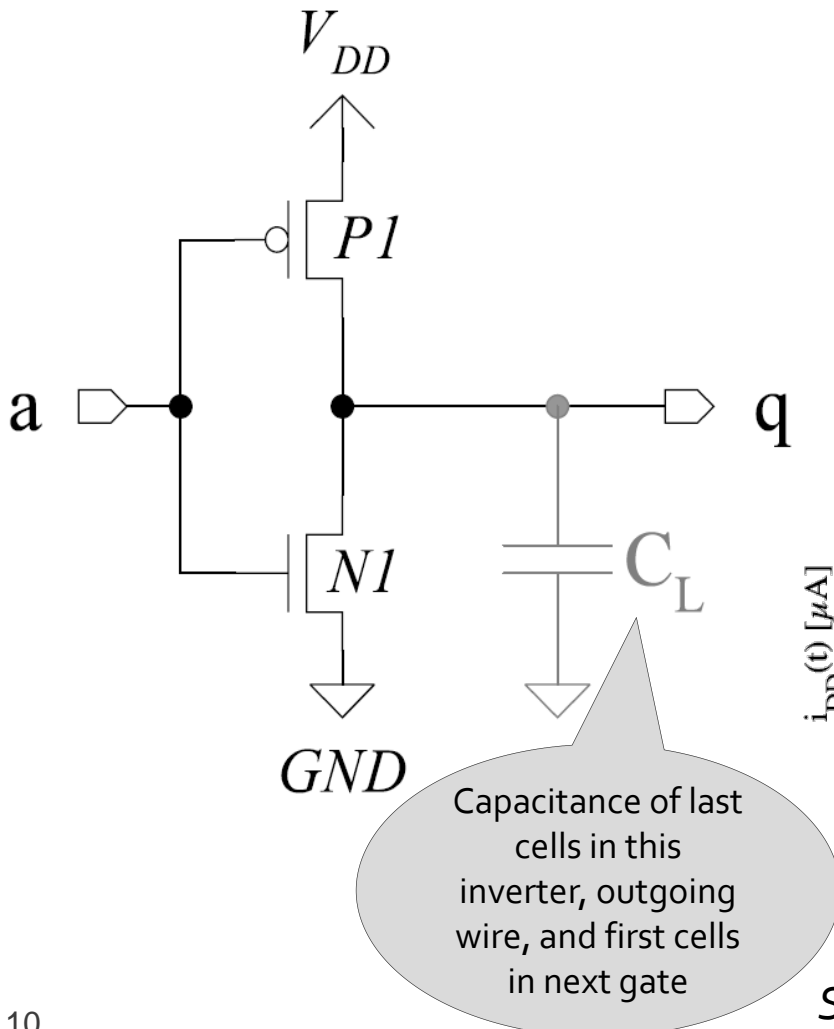
Data-dependent power consumption: gate capacitance



- The power consumption of a CMOS gate depends on the data:
 - q: 0→0 almost no power consumption
 - q: 1→1 almost no power consumption
 - q: 0→1 high power consumption (proportional to C_2)
 - q: 1→0 high power consumption (proportional to C_1)



Data-dependent power consumption: wire capacitance



Source: DPA Book

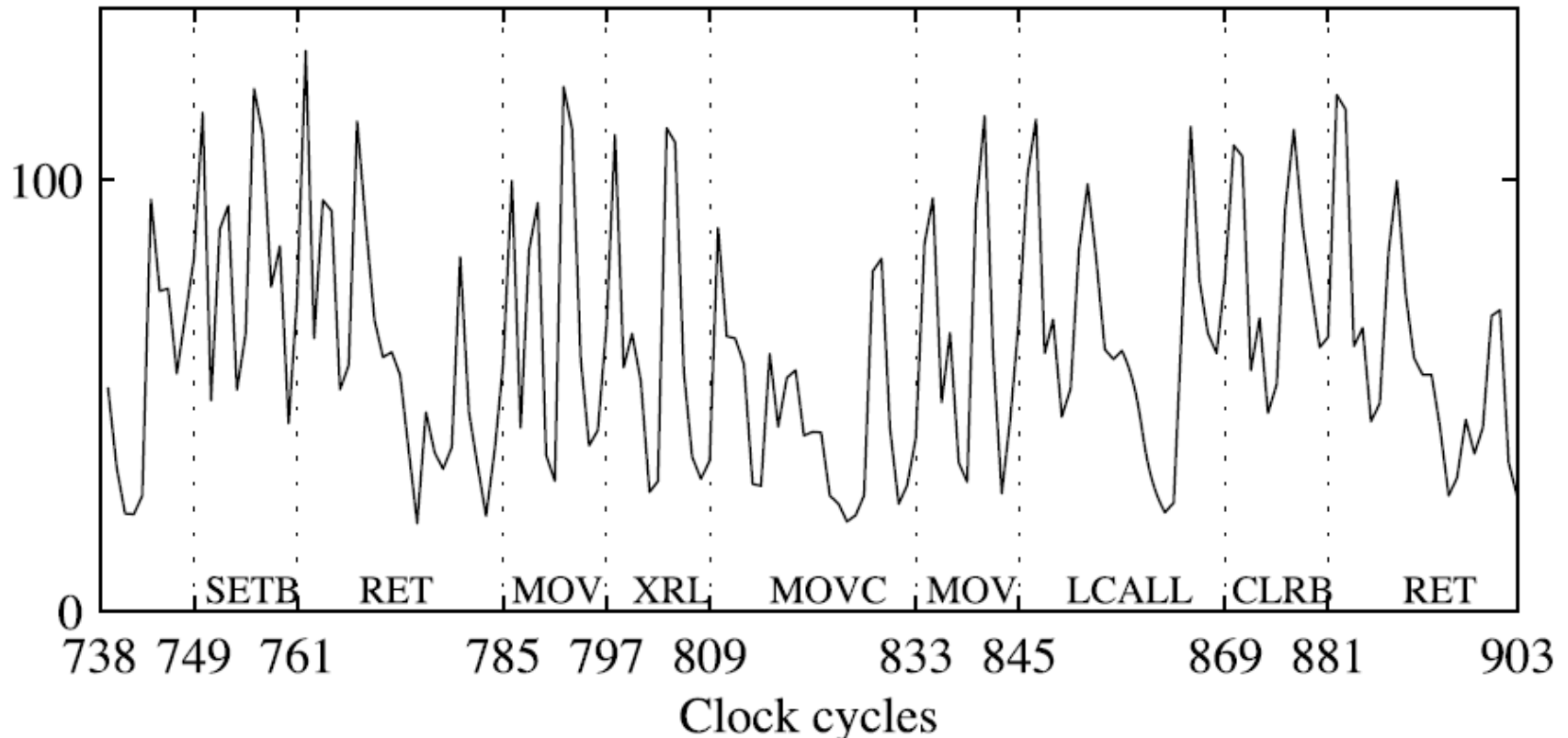


Power Depends on Instruction, Locally

```
LCALL SET_ROUND_TRIGGER
MOV A,ASM_input + 0      ; load a0
XRL A,ASM_key + 0       ; add k0
MOVC A,@A + DPTR        ; S-box look-up
MOV ASM_input, A        ; store a0
LCALL CLEAR_ROUND_TRIGGER
```



Power Depends on Instruction, Locally



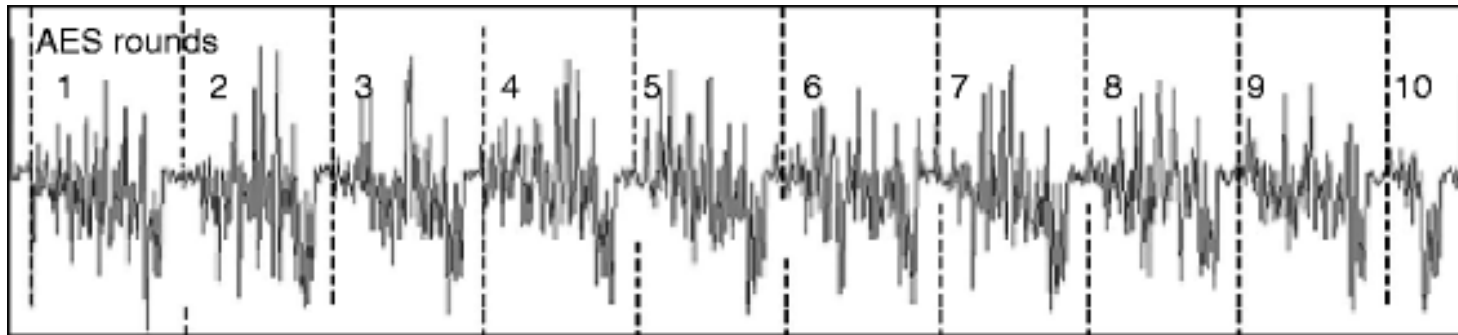
Source: DPA Book



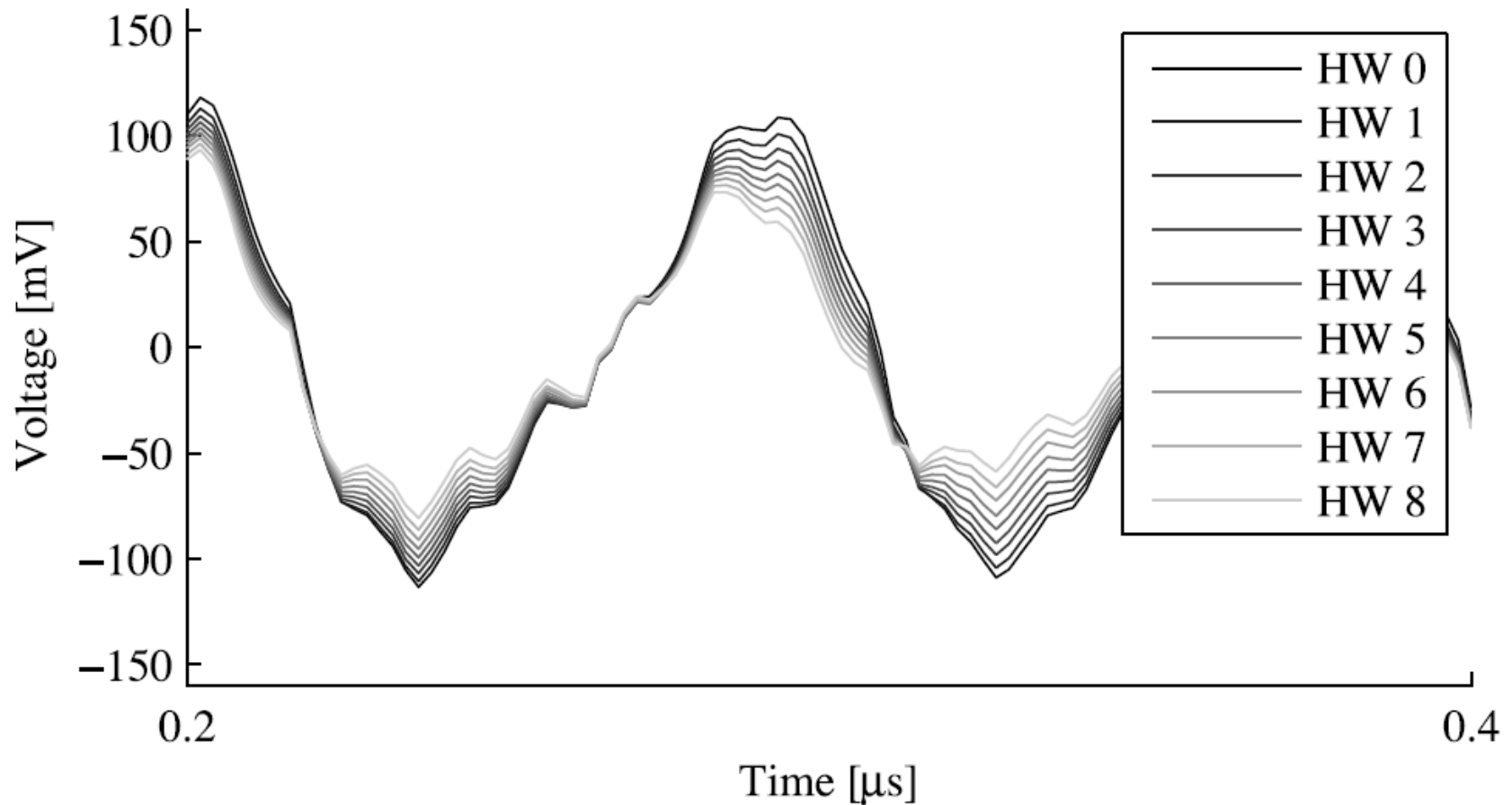
Power Depends on Instruction, Globally

Example: AES

[Joye Olivier 2011]



Power Depends on Data



Source: DPA Book



Power Analysis

- Simple Power Analysis
- Differential Power Analysis
 - Warm-up Correlation Power Analysis
 - Full Correlation Power Analysis



Power Analysis Attack Scenario

- Plaintexts and ciphertexts may be **chosen**, **known** or **unknown**



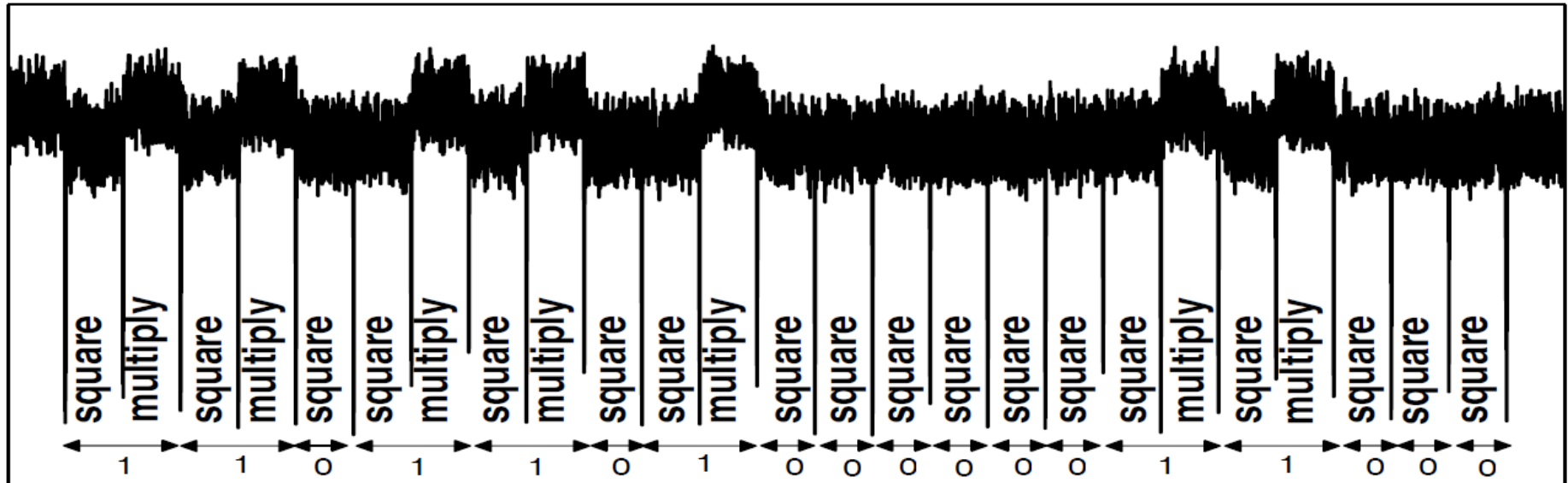
Theory of power analysis

- Power consumption is **variable**
- Power consumption depends on **instruction**
- Power consumption depends on **data**



Simple Power Analysis (SPA)

Example: square-and-multiply RSA exponentiation.



- Pros:
 - Single trace (or average of a few) may suffice
- Cons:
 - Detailed reverse engineering
 - Long manual part
 - Hard to handle bad signal-to-noise ratio, especially for small events, (e.g., AES)

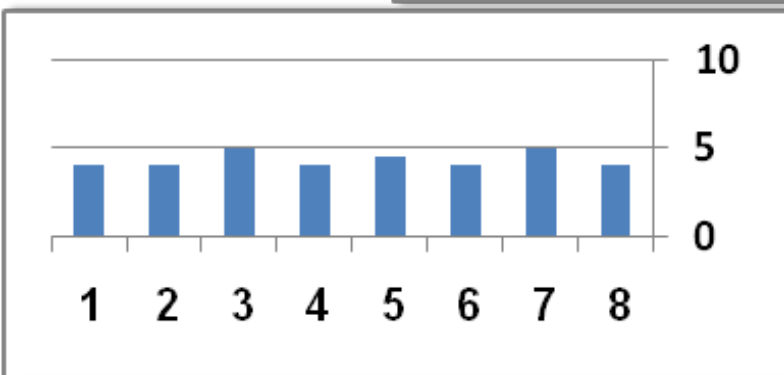
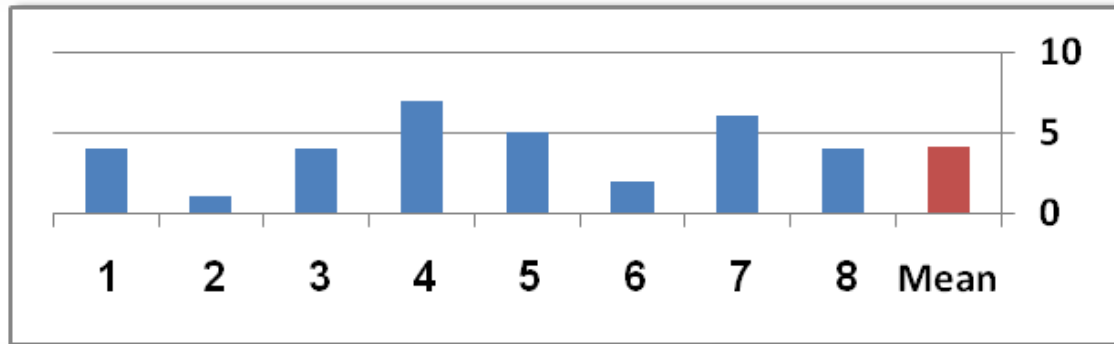


Differential Power Analysis (DPA)

- Use statistical properties of traces to recover key
- Pros:
 - Very limited reverse engineering
 - Harder to confuse
- Cons:
 - Large amount of traces
- Two main types of DPA:
 - Difference of means (traditional DPA)
 - **Correlation power analysis (CPA)**

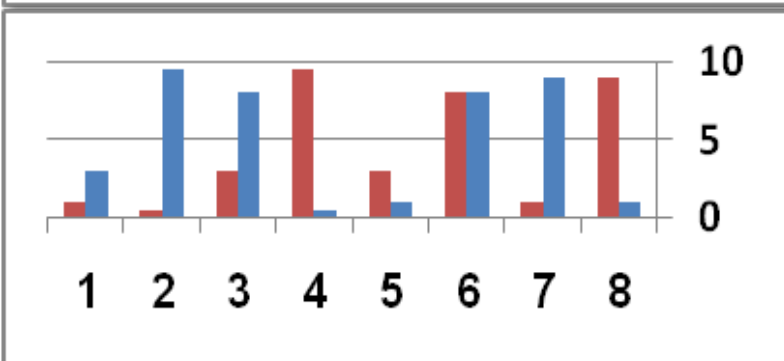
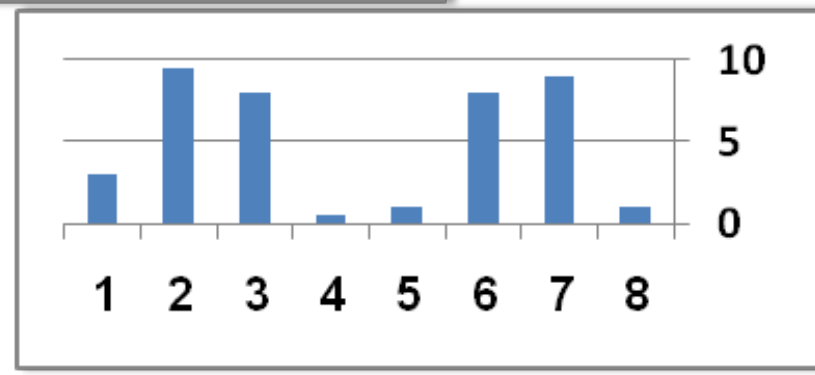


Mean, variance, correlation



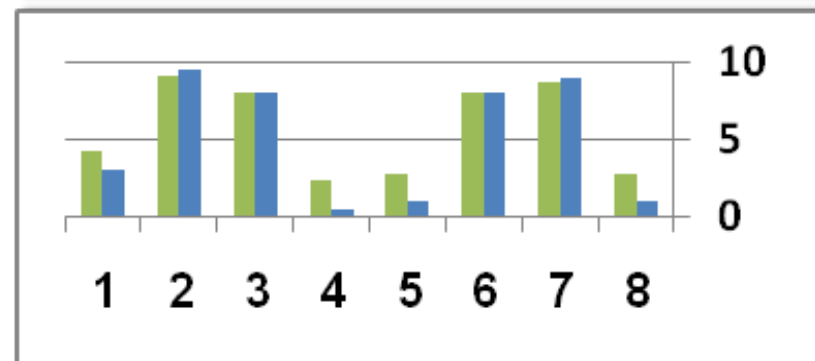
⇐ Low
Variance

High ⇒
Variance



⇐ Low
Correlation

High ⇒
Correlation



CPA Basics

- We want to discover the **correct key value** (c_k) and **when it is used** (c_t)
- Idea:
 - On the **correct time**, the power consumption of **all traces** is **correlated** with the **correct key**
 - On **other times** and **other keys** the traces should show **low correlation**



Warm-up CPA

- Assume **plaintext** and **correct key** are known but **correct time** is unknown
- Form **hypothesis** and test it
- Good **hypothesis function** $f(p, k)$
 - Depends on **known plaintext** and
 - Depends on **small amount of key bits**
 - **Deterministic**
 - Sufficiently random (e.g., non-linear), sensitive to small changes in p and k
- Maps to power consumption using a **power consumption model**

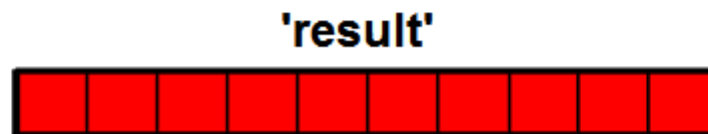
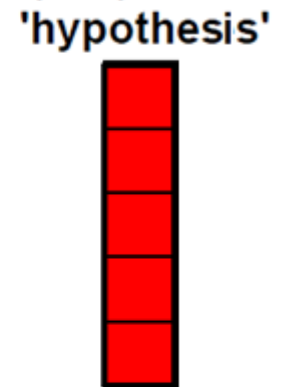
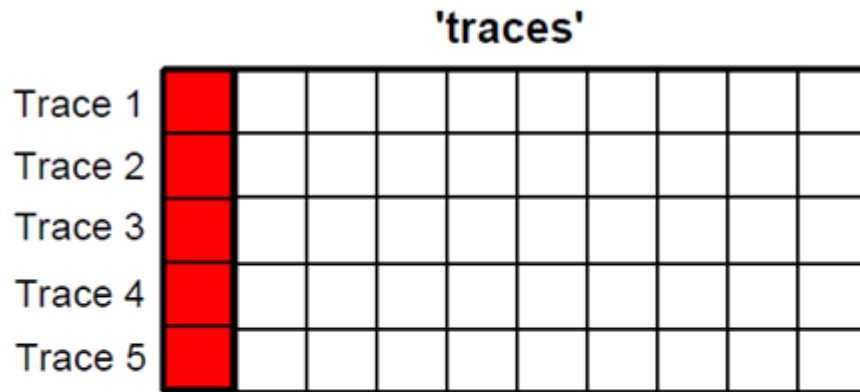


Warm-up CPA in Numbers

- **1000 traces**, each consisting of **1 million points**
- Each trace uses a different known plaintext – **1000 plaintexts**
- **1 known key**
- Hypothesis is vector of **1000 hypothetical power values**
- Output of warm-up CPA: vector of **1 million correlation values** with peak at c_t



Warm-up CPA in Pictures



Full CPA

- Plaintext is **known**, but **correct key** and **correct time** unknown
- Idea: run warm-up CPA many times in parallel
- Create **many competing hypotheses**



Full CPA in Numbers

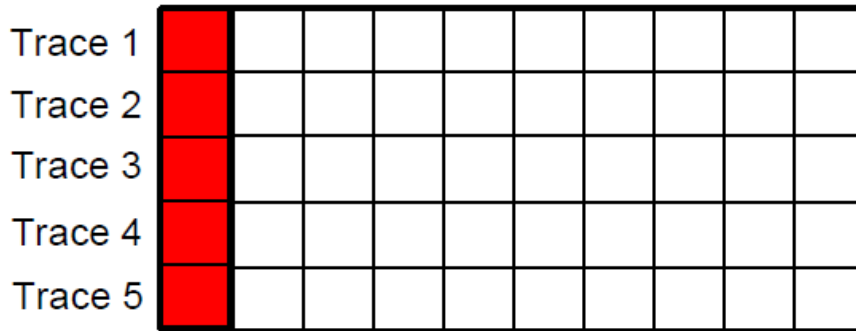
- **1000 traces**, each consisting of **1 million points**
- Each trace uses a different known plaintext – **1000 plaintexts**
- Key is unknown – **256 guesses** for first byte
- Hypothesis is matrix of **1000X256 hypothetical power values**
- Output of full CPA: matrix of **1,000,000X256 correlation values** with peak at (c_k, c_t)



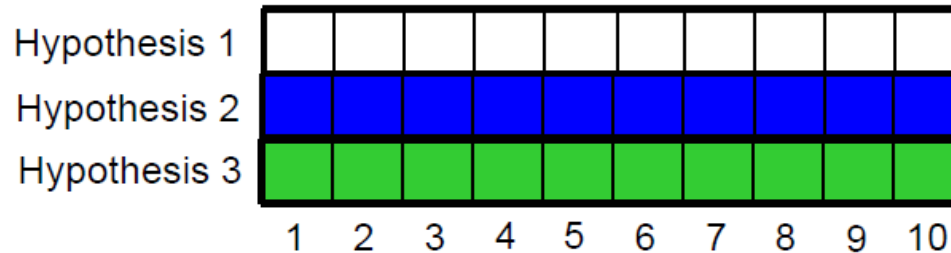
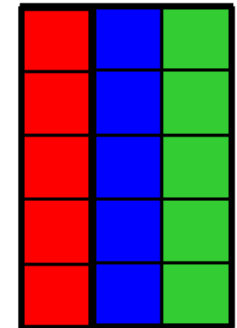
Full CPA in Pictures

`corr_traces = result(dpa_obj);`

(5x10) matrix 'traces'



(5x3) matrix 'hypotheses'



(3x10) matrix 'corr_traces'



Full attack

- Acquire traces
- For every key fragment and corresponding hypothesis function:
 - Compute matrix of hypotheses
 - Compute correlation (score) matrix
 - Find maximum in correlation matrix and deduce value of key fragment



Discussion: effects and tradeoffs of parameters

- Number of traces
- Trace length
- Number of hypotheses (i.e., number of relevant key bits affecting the hidden state being modeled)



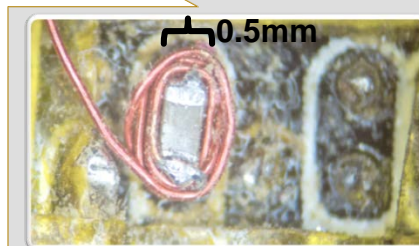
Assumptions/complications/mitigation

ASSUMPTIONS/BOTTLENECKS

- Alignment (shifts, fragmentation)
 - Handling algorithmically (e.g., cross correlation, string alignment)
 - Physical synchronization via triggering
- Data size: #traces, trace length
- Assuming fixed over all traces:
 - Key
 - Computation progress and flow
- Measurement quality
 - Physical access, noise, transmission
 - Equipment cost
 - Expertise

MITIGATIONS

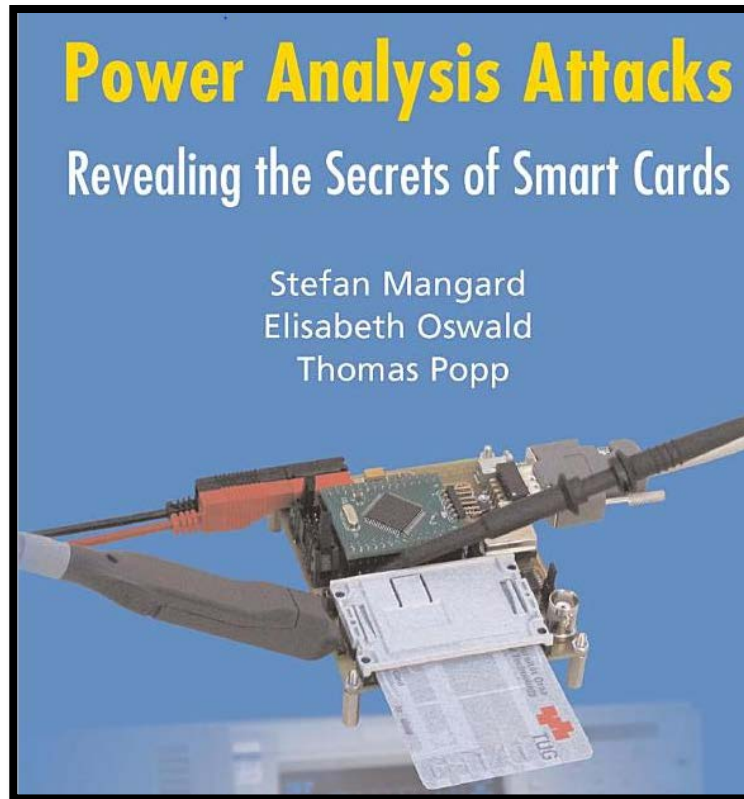
- Timing
 - Unstable clock
 - Randomize control/instruction flow
 - Insert random/key/plaintext-dependent delays
- Change protocol to roll keys often
- Add power noise



Example: probing an “decoupling capacitor” (SMD 0402) close to the microcontroller chip.
[O’Flynn 2013]



Further Reading



<http://www.dpabook.org>

<http://www.springerlink.com/content/go1q1k>

